Self-Stabilizing Silent Disjunction in an Anonymous Network

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ICDCN 2013, 04/01/2012, Mumbai





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- 2 Basic Principle: Flooding
- 3 Color Waves to prevent Endless Loops
- Use Energy to prove Stabilization







- 2 Basic Principle: Flooding
- 3 Color Waves to prevent Endless Loops
- 4 Use Energy to prove Stabilization
- 5 Conclusion



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Self-Stabilization



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System Configurations



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Illegitimate Configurations

Legitimate Configurations



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Tolerance to Transient Faults



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• Anonymous Network of Processes.



- Anonymous Network of Processes.
 - Each Process has Input Bit.



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- Anonymous Network of Processes.
 - Each Process has Input Bit.
 - **Double Circle** \iff Input Bit = 1. Others are 0.



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 - Output Bit of each Process is Disjunction of all Input Bits of Processes in its Component.



- Anonymous network of Processes.
 - Each Process has Input Bit.
 - Double Circle \iff Input Bit = 1. Others are 0.
 - Output Bit shown inside circle.
 - Output Bit of each Process is Disjunction of all Input Bits of Processes in its Component.
- (Construct BFS Forest Rooted at Processes with Input 1.)







3 Color Waves to prevent Endless Loops

4 Use Energy to prove Stabilization

5 Conclusion



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Distributed Disjunction

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• Works if Clean Start.



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- Works if Clean Start.
- If Arbitrary Start:



- Works if Clean Start.
- If Arbitrary Start:
 - Works if Some Process has Input 1.



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- Works if Clean Start.
- If Arbitrary Start:
 - Works if Some Process has Input 1.
 - All Processes have Input 0: might go into Endless Loop!



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• All Input Bits 0.



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- All Input Bits 0.
- Clean Configuration is Final.



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• Some Input Bits = 1.



- Some Input Bits = 1.
- Clean Configuration is Not Final.



- Some Input Bits = 1.
- Clean Configuration is not Final.
- Clusterheads Execute Reset.



- Some Input Bits = 1.
- Clean Configuration is not Final.
- Clusterheads Execute Reset.
- Adjacent Processes Execute Join.



- Some Input Bits = 1.
- Clean Configuration is not Final.
- Clusterheads Execute Reset.
- Adjacent Processes Execute Join.



- Some Input Bits = 1.
- Clean Configuration is not Final.
- Clusterheads Execute Reset.
- Adjacent Processes Execute Join.
- Flooding: All Parent Pointers and Levels are Computed in at Most (1 + Diam) Rounds.





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• Arbitrary Initial Configuration?



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- Arbitrary Initial Configuration.
- No Problem if Some Process has Input 1.



- Arbitrary Initial Configuration.
- No Problem if Some Process has Input 1.
- Otherwise, might go into Endless Loop.




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Distributed Disjunction

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• Some Input Bits 1.



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- Some Input Bits 1.
- Red = Enabled to Reset.
- Green = Enabled to Join.



- Some Input Bits 1.
- Red = Enabled to Reset.
- Green = Enabled to Join.
- Clusterheads Reset.



- Some Input Bits 1.
- Red = Enabled to Reset.
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- Clusterheads Reset.
- Flooding Moves at Speed 1.



- Some Input Bits 1.
- Red = Enabled to Reset.
- Green = Enabled to Join.
- Clusterheads Reset.
- Flooding Moves at Speed 1.



- Some Input Bits 1.
- Red = Enabled to Reset.
- Green = Enabled to Join.
- Clusterheads Reset.
- Flooding Moves at Speed 1.
- Convergence within 1 + Diam Rounds.





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• All Input Bits 0.



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- All Input Bits 0.
- Output Bits Inside Circles.





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- All Input Bits 0.
- Output Bits Inside Circles.
- Red = Enabled to Reset.





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- All Input Bits 0.
- Output Bits Inside Circles.
- Red = Enabled to Reset.
- Green = Enabled to Join.





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- All Input Bits 0.
- Output Bits Inside Circles.
- Red = Enabled to Reset.
- Green = Enabled to Join.
- Chain Deletes at Head End.
- But Recruits at Tail (Leaf) End.





- All Input Bits 0.
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- Red = Enabled to Reset.
- Green = Enabled to Join.
- Chain Deletes at Head End.
- But Recruits at Tail (Leaf) End.
- Keeps Going Around!





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- All Input Bits 0.
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- Red = Enabled to Reset.
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- Chain Deletes at Head End.
- But Recruits at Tail (Leaf) End.
- Keeps Going Around.
- Return to First Configuration, Except for Levels.





- All Input Bits 0.
- Output Bits Inside Circles.
- Red = Enabled to Reset.
- Green = Enabled to Join.
- Chain Deletes at Head End.
- But Recruits at Tail (Leaf) End.
- Keeps Going Around.
- Return to First Configuration, Except for Levels.
- Endless!







2 Basic Principle: Flooding

3 Color Waves to prevent Endless Loops

4 Use Energy to prove Stabilization

5 Conclusion



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• Prevent Indefinite Growth of Fictitious Trees.



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• Prevent Indefinite Growth of Fictitious Trees.

Side Effects of Color Waves



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• Prevent Indefinite Growth of Fictitious Trees.

Side Effects of Color Waves

• Slow Down Algorithm by a Factor of Three



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• Prevent Indefinite Growth of Fictitious Trees.

Side Effects of Color Waves

- Slow Down Algorithm by a Factor of Three
- After Legitimacy, Color Waves could Run Forever.



• Prevent Indefinite Growth of Fictitious Trees.

Side Effects of Color Waves

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Counteract Effect with Done Waves



• Prevent Indefinite Growth of Fictitious Trees.

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Counteract Effect with Done Waves

• Convergecast: Leaves Detect Algorithm Done



• Prevent Indefinite Growth of Fictitious Trees.

Side Effects of Color Waves

- Slow Down Algorithm by a Factor of Three
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Counteract Effect with Done Waves

- Convergecast: Leaves Detect Algorithm Done
- Root (Clusterhead) Color Freezes.



• Prevent Indefinite Growth of Fictitious Trees.

Side Effects of Color Waves

- Slow Down Algorithm by a Factor of Three
- After Legitimacy, Color Waves could Run Forever.

Counteract Effect with Done Waves

- Convergecast: Leaves Detect Algorithm Done
- Root (Clusterhead) Color Freezes.
- Color Lock Results Within O(Diam) Rounds.



• Prevent Indefinite Growth of Fictitious Trees.

Side Effects of Color Waves

- Slow Down Algorithm by a Factor of Three
- After Legitimacy, Color Waves could Run Forever.

Counteract Effect with Done Waves

- Convergecast: Leaves Detect Algorithm Done
- Root (Clusterhead) Color Freezes.
- Color Lock Results Within O(Diam) Rounds.

• Silence.





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• Process with Output Bit = 1 has Color: 0 = 1, 1 = 1.



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- Process with Output Bit = 1 has Color: 0 = (1, 1 = (1, 1)).
- If Process X Executes Join, Attaching to Process Y:



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- Process with Output Bit = 1 has Color: 0 = (1, 1 = (1, 1)).
- If Process X Executes Join, Attaching to Process Y:
 - Y must have Color 1: $\bigcup_{x} \bigcup_{y} \bigcup_{y}$



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- Process with Output Bit = 1 has Color: 0 = 0, 1 = 0.
- If Process X Executes Join, Attaching to Process Y:
 - Y must have Color 1: O
 Color of X Becomes 0: O



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- Process with Output Bit = 1 has Color: 0 = (1, 1 = (1, 1)).
- If Process X Executes Join, Attaching to Process Y:
- Process X can Change Color if the Following Conditions Hold:


- Process with Output Bit = 1 has Color: 0 = (1, 1 = (1, 1)).
- If Process X Executes Join, Attaching to Process Y:
 - Y must have Color 1: O 0 X Y
 Color of X Becomes 0: Y Y

• Process X can Change Color if the Following Conditions Hold:

• X has Same Color as its Parent, or is Clusterhead:

$$\begin{array}{c} 1 \\ X \\ X \end{array} \rightarrow \begin{array}{c} 1 \\ X \\ X \end{array} \text{ or } \begin{array}{c} 1 \\ X \\ X \end{array} \rightarrow \begin{array}{c} 1 \\ X \\ X \end{array} \text{ or } \begin{array}{c} 1 \\ X \\ X \end{array}$$



- Process with Output Bit = 1 has Color: 0 = (1, 1 = (1, 1)).
- If Process X Executes Join, Attaching to Process Y:
 - Y must have Color 1: X must have Color

• Process X can Change Color if the Following Conditions Hold:

• X has Same Color as its Parent, or is Clusterhead:

 $(1 \rightarrow 1)$ or $(1 \rightarrow 1)$ or (1) or (1) or (1) or



or ⁽¹⁾

- Process with Output Bit = 1 has Color: 0 = (1, 1 = (1, 1)).
- If Process X Executes Join, Attaching to Process Y:

• Process X can Change Color if the Following Conditions Hold:

• X has Same Color as its Parent, or is Clusterhead:

(1) or (1) \rightarrow (1) or (1) χ

- Children have Opposite Color:
- If Color = 1, X Cannot Change Color if Any Neighbor is

Enabled to Attach to X:



or

or

- Process with Output Bit = 1 has Color: 0 = (1, 1 = (1, 1)).
- If Process X Executes Join, Attaching to Process Y:
 - Y must have Color 1: O_x
 - Color of **X** Becomes 0:

• Process X can Change Color if the Following Conditions Hold:

• X has Same Color as its Parent, or is Clusterhead:

 $1 \rightarrow 1$ or $1 \rightarrow 1$ or $1 \rightarrow 1$ or $1 \rightarrow 1$ or $1 \rightarrow 1$

- If Color = 1: X Cannot Change Color if Any Neighbor is Enabled to Attach to X:
- When Clusterhead Changes Color, a Color Wave is Absorbed

or

- Process with Output Bit = 1 has Color: 0 = (1, 1 = (1, 1)).
- If Process X Executes Join, Attaching to Process Y:
 - Y must have Color 1: O_{X}
 - Color of **X** Becomes 0:

• Process X can Change Color if the Following Conditions Hold:

• X has Same Color as its Parent, or is Clusterhead:

 $(1)_{\chi} \rightarrow (1)_{\chi}$ or $(1)_{\chi} \rightarrow (1)_{\chi}$ or $(1)_{\chi}$ or

- If Color = 1: X Cannot Change Color if Any Neighbor is Enabled to Attach to X:
- When Clusterhead Changes Color, a Color Wave is Absorbed
 False Roots Cannot Absorb Color Waves.

or



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Distributed Disjunction

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• Chain Example.

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- Chain Example.
- One Process has Input = 1.

- Chain Example.
- One Process has Input = 1.
- Arrow Shows Flow of Time.

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- Chain Example.
- One Process has Input = 1.
- Arrow Shows Flow of Time.
- Color Waves Move
 Toward Clusterhead

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- Chain Example.
- One Process has Input = 1.
- Arrow Shows Flow of Time.
- Color Waves Move Toward Clusterhead
- Growth Rate = 1/3.

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- Chain Example.
- One Process has Input = 1.
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- Color Waves Move Toward Clusterhead
- Growth Rate = 1/3.

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- Consider a False Root, R.
- **R** is Enabled Only to **Reset**, and thus Cannot Change Color.
- Eventually Color Lock.
- Tree Rooted at **R** Cannot Grow Forever.
- How can you Prove That?
- Use Energy!



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- 2 Basic Principle: Flooding
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5 Conclusion







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Distributed Disjunction

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• Energy(X): Positive Integer for X of Output = 1.



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- Energy(X): Positive Integer for X of Output = 1.
- Defined Recursively.



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• Energy(X) = 1 if X is Leaf of Color = 0.



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- Energy(X): Positive Integer for X of Output = 1.
- Defined Recursively.
 - Energy(X) = 1 if X is Leaf of Color = 0.
 - Energy(X) = 2 if X is Leaf of Color = 1.



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- Energy(X): Positive Integer for X of Output = 1.
- Defined Recursively.
 - Energy(X) = 1 if X is Leaf of Color = 0.
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 - X Not Leaf: Energy(X) = Maximum of:



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 - Energy(X) = 1 if X is Leaf of Color = 0.
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 - X Not Leaf: Energy(X) = Maximum of:
 - 1 + Energy of any Child of Opposite Color.



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- Theorem: No Action of a Process of Input = 0 can Increase Maximum Energy of the Network.



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 - Energy(X) = 1 if X is Leaf of Color = 0.
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 - X Not Leaf: Energy(X) = Maximum of:
 - 1 + Energy of any Child of Opposite Color.
 - 2 + Energy of any Child of Matching Color.
- Theorem: No Action of a Process of Input = 0 can Increase Maximum Energy of the Network.
- Theorem: If All Processes have Input = 0: Maximum Energy of the Network Decreases every Round. Hence Convergence After O(n) Rounds.



- Energy(X): Positive Integer for X of Output = 1.
- Defined Recursively.
 - Energy(X) = 1 if X is Leaf of Color = 0.
 - Energy(X) = 2 if X is Leaf of Color = 1.
 - X Not Leaf: Energy(X) = Maximum of:
 - 1 + Energy of any Child of Opposite Color.
 - 2 + Energy of any Child of Matching Color.
- Theorem: No Action of a Process of Input = 0 can Increase Maximum Energy of the Network.
- Theorem: If All Processes have Input = 0: Maximum Energy of the Network Decreases every Round, Hence Convergence After O(n) Rounds.





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• Legitimate Configuration After O(n) rounds



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- Legitimate Configuration After O(n) rounds
- How do we Stop Color Waves from Continuing Forever?



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- Done Waves:



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- How do we Stop Color Waves from Continuing Forever?
- Done Waves:
 - Leaf Initiates when it Detects (Local) Legitimacy.



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Silence

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Silence

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- Clusterhead (That is, Process with Input = 1) Becomes Color Frozen. Cannot Change Color.



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Silence

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- Done Wave Convergecast.
- Clusterhead (That is, Process with Input = 1) Becomes Color Frozen. Cannot Change Color.
- Color Lock: Within O(Diam) Rounds: Silence is Achieved.



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Asynchronous Example Computation



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- Asynchronous Example Computation
- All Input Bits 0.



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- All Input Bits 0.
- Red = Enabled to Reset.



- Asynchronous Example Computation
- All Input Bits 0.
- Red = Enabled to Reset.
- Green = Enabled to Join.



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- All Input Bits 0.
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- Maximum **Energy** Initially = 10.



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- Configuration is Now Legitimate.







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- Blue = Enabled to Change Color.
- Maximum Energy Initially = 10.
- Maximum Energy Decreases each Round.
- No Further Growth Possible.
- Configuration is Now Legitimate.
- Configuration is Now Final. Silent.









- 2 Basic Principle: Flooding
- 3 Color Waves to prevent Endless Loops
- 4 Use Energy to prove Stabilization





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• We introduced the Disjunction Problem

• Application: Clustering

Solution:

- Locally Shared Memory Model (Unfair Daemon)
- O(log Diam + log △) space per process
- Stabilization Time: O(n) rounds

• Stabilization Time in O(Diam) rounds ?

- Should be hard!
- As difficult as Self-Stabilizing Leader Election in identified networks in O(Diam) rounds ...



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Thank You!



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