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Verification of analog and mixed-signal circuits using hybrid systems techniques

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Plan

1. Introduction
2. Verification of analog circuits
3. Verification of mixed-signal circuits
4. Conclusion

Introduction

- Analog and mixed-signal circuits: increasing applications
- Traditional design of analog ICs: mixture of expertise, lots of manual calculations
- **Numerical simulation** not suitable to prove that the circuit satisfies the specification under all possible inputs and parameter values
- **Symbolic analysis methods**
 - Most tools are developed for linear or linearized circuits
 - For nonlinear circuits: reduction to linear or weakly nonlinear systems, limited by accuracy trade-off
 - Time-domain symbolic analysis: less well-developed

Our goal and approach

Goal: We want to analyze

- *analog* and *mixed-signal* circuits whose continuous dynamics are described by a system of differential algebraic equations (DAEs)
- *time-domain properties*
- under *all* possible *input signals* and *parameter values*

Approach: use the verification techniques for continuous/hybrid systems

- Extend the reachability analysis techniques for ordinary differential equations (ODEs) to DAEs
- Optimization-based verification methods (to find the worst case behaviors)

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Model of analog circuits

Behavior of an analog circuit can be described by **nonlinear differential algebraic equations** (DAE):

$$F(x(t), \dot{x}(t), y(t), u(t), p) = 0$$

u : inputs; x : state variables (internal voltages and currents); y : outputs;
 p : circuit parameters

Analyze time-domain properties

- Characterize solutions under all possible input signals and parameters
⇒ *compute reachable sets*
 - Reachability analysis techniques for *ordinary differential equations* developed in the context of *hybrid systems* (e.g. tools CheckMate (CMU), d/dt (Verimag), level set methods (Stanford))
- ⇒ **Extend ODE reachability analysis techniques to DAEs**

Reachability analysis of DAEs

$$F(x, \dot{x}, y, u, p) = 0$$

- Differential index: minimal number of differentiations required to solve for the derivatives of variables
- Practical electronic circuits: often index 1 or 2.
- We focus on DAEs of index 1

Reachability analysis of DAEs (cont'd)

We study the equivalent semi-explicit form of the DAE system:

$$\begin{aligned}\dot{x} &= f(x, y, p) \\ 0 &= g(x, y, p)\end{aligned}$$

- **Transforming into ODEs :**

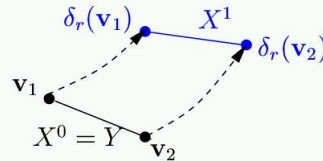
Differentiating the algebraic eq. once gives $\dot{y} = -g_y^{-1}g_x f$ where $g_y(x, y) = \partial g / \partial y$. (Note that the DAEs are of index 1)
 \Rightarrow Obtain augmented ODEs with variables $z = (x, y)^T$:

$$\dot{z} = (f, -g_y^{-1}g_x f)^T = \tilde{f}$$

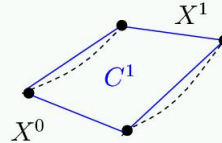
- **Retain the algebraic constraint** and interpret the original DAEs as the augmented **ODEs on a manifold** :

$$\begin{aligned}\dot{z} &= \tilde{f}(z, p) \\ 0 &= g(z, p)\end{aligned}$$

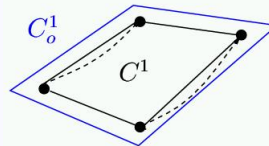
ODE reachability technique of d/dt



$$X^1 = \text{conv}\{\delta_r(v_1), \delta_r(v_2)\}$$



$$C^1 = \text{conv}(X^0 \cup X^1)$$



$$C_o^1 = \text{bloat}(C^1, \varepsilon)$$

Linear system $\dot{x}(t) = Ax(t) + u(t)$

Second order method (approximation error is of order $\mathcal{O}(r^2)$ where r is time step) [AsarinDangMaler01]

ODE reachability technique of d/dt (cont'd)

Extension to **nonlinear systems** : $\dot{x}(t) = f(x(t))$ [AsarinDangGirard03]

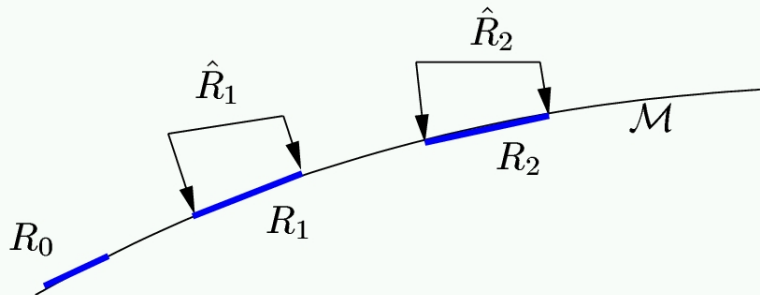
- Approximation of f by a piecewise affine function over a simplicial mesh
- Approximation error μ modeled as uncertain input: $\dot{x}(t) = A_i x(t) + u(t)$ where $\|u\| \leq \mu \Rightarrow$ conservativeness
- Second order method (if f is a C^2 function)

Reachability analysis of ODEs on a manifold

We study a system of ODEs on a manifold:

$$\begin{aligned}\dot{z}(t) &= f(z(t)) \\ 0 &= g(z(t)) \Rightarrow \text{defining a manifold } \mathcal{M} \\ z(0) &\in R_0\end{aligned}$$

Combining reachability computations techniques for ODEs and ideas from *geometric integration using projection* [Lubich, Hairer, Wanner 2003]



Reachability algorithm for ODEs on a manifold

R_0 : initial set
repeat $k = 0, 1, \dots$
 $\hat{R}_{k+1} = \text{Reach}_{[0,r]}(R_k)$ /* computed for the augmented ODEs */
 $R_{k+1} = \Pi_{\mathcal{M}}(\hat{R}_{k+1})$ /* project on the manifold \mathcal{M} */
until $R_{k+1} = \bigcup_{i=1}^k R_i$

- Projection:

$$\Pi_{\mathcal{M}}(\hat{z}) = \arg \min_z |\hat{z} - z| \quad \text{subject to } g(z) = 0$$

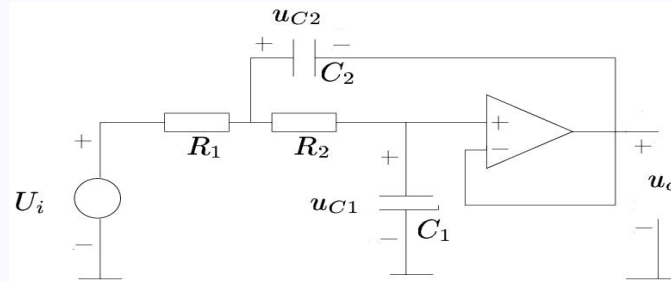
- **Convergence** : same order as the convergence order of the technique for ODEs (projection does not deteriorate the convergence)
- Second order method

Computational issues

- Nonlinearities in analog circuits often restricted to the device models.
- Not all the nonlinearities are important for the property under study.
- ‘Hybridization’
 - approximation of a complex function by a piecewise simpler one
 - can be used for complexity reduction
 - gives rise to a hybrid system
- Hybrid systems: good model of many circuit element characteristics (which are often smoothed to avoid numerical instability)

Example: Biquad low pass filter

[Hartong,Hedrich,Barke 2002]



$$\dot{u}_{C1} = \frac{u_{C2} + u_o - u_{C1}}{C_1 R_2} \quad \dot{u}_{C2} = \frac{U_i - u_{C2} - u_o}{C_2 R_1} - \frac{u_{C2} + u_o - u_{C1}}{C_2 R_2}, \quad (1)$$

$$u_o - V_{max} \tanh\left(\frac{(u_{C2} - u_o)V_e}{V_{max}}\right) + U_{om} = 0, \quad (2)$$

$$U_{om} = \mathcal{V}(i_o), \quad i_o = -C_2 \dot{u}_{C2}, \quad (3)$$

$$\mathcal{V}(i_o) = K_1 i_o + 0.5 \sqrt{K_1 i_o^2 - 2K_2 i_o I_s + K_1 I_s^2 + K_2} - 0.5 \sqrt{K_1 i_o^2 + 2K_2 i_o I_s + K_1 I_s^2 + K_2}. \quad (4)$$

Differentiating (2) \Rightarrow nonlinear ODEs on a manifold with variables $\mathbf{z} = (u_{C1}, u_{C2}, u_o)$

Biquad low pass filter: Approximation by hybridization

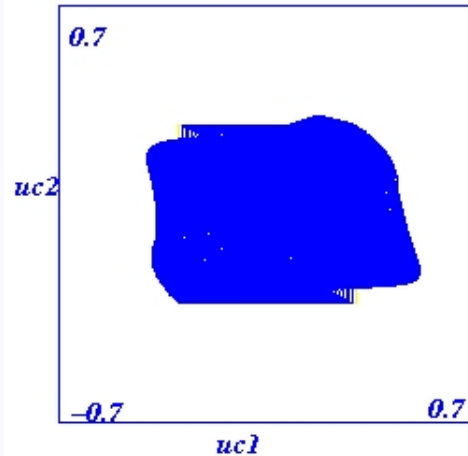
- *tanh* approximated by polynomial or piecewise affine \Rightarrow loss of important behaviors of the circuit
- the nonlinear characteristics $U_{om} = \mathcal{V}(i_o)$ can be approximated by a piecewise-affine function:

$$U_{om} = \begin{cases} 1e10 i_o + 0.5e8 & \text{if } i_o \leq -i_s \\ 0 & \text{if } -i_s < i_o < i_s \\ 1e10 i_o - 0.5e8 & \text{if } i_o \geq i_s \end{cases} \quad (5)$$

\Rightarrow Hybrid automaton with 3 continuous modes

Biquad low pass filter: verification results

The property to verify is the *absence of overshoots*.



- $C_1 = 0.5e - 8$, $C_2 = 2e - 8$, and $R_1 = R_2 = 1e6$ (highly damped case)
- The initial set: $u_{C1} \in [-0.3, 0.3]$, $u_{C2} \in [-0.3, 0.3]$ and $u_o \in [-0.2, 0.2]$

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Verification of mixed-signal circuits

Mixed-signal circuits are circuits mixing analog signals (real-valued) and digital signals (boolean-valued).

Model: We consider circuits that can be modeled by:

$$F(x(k), x(k+1), u(k), \delta(k), p) = 0, \quad k \in \mathbb{N},$$

- $\delta(k) \in \{0, 1\}^s$ (modelling logical part)
- $u(\cdot) = (u(k))_{k \in \mathbb{N}}$ input sequence
- $x(k)$ state at time k , uniquely defined from $x(0)$ and $u(\cdot)$

Safety verification problem: given a safe set \mathcal{S} , verify that the system always remains in \mathcal{S} over a bounded horizon N

Optimisation based verification: Find the *worst-case* behavior and check if it satisfies the property

Optimisation based verification

- Define an objective function J such that:

$$J(x) \text{ is positive} \Leftrightarrow x \text{ is outside } \mathcal{S}$$

- To prove that *the system is safe at time k* : checking whether **the maximum value of $J(x(k))$ is not positive** by solving:

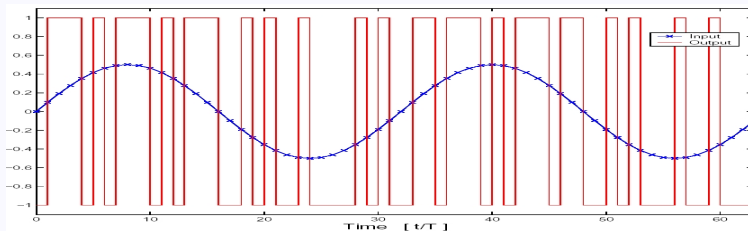
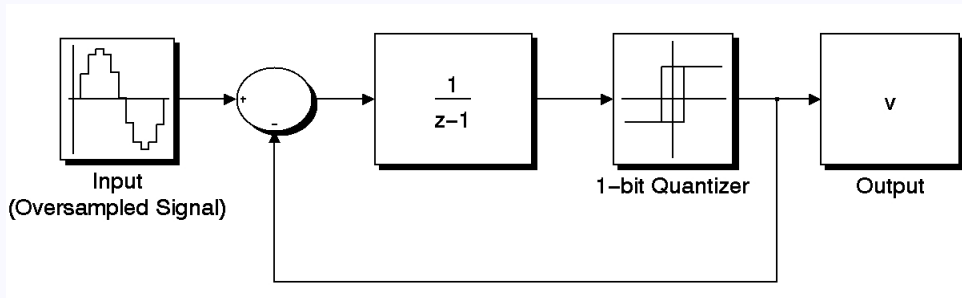
$$\begin{aligned} & \max J(x(k)), \\ & \text{s.t. } F(x(t), x(t+1), u(t), \delta(t), p) = 0, \\ & \quad u(t) \in \mathcal{U}, \delta(t) \in \{0, 1\}^s, t \in \{0, 1, \dots, k-1\}, \\ & \quad x(0) \in X_0. \end{aligned}$$

This constrained optimisation problem is mixed since $\delta(t)$ are integers

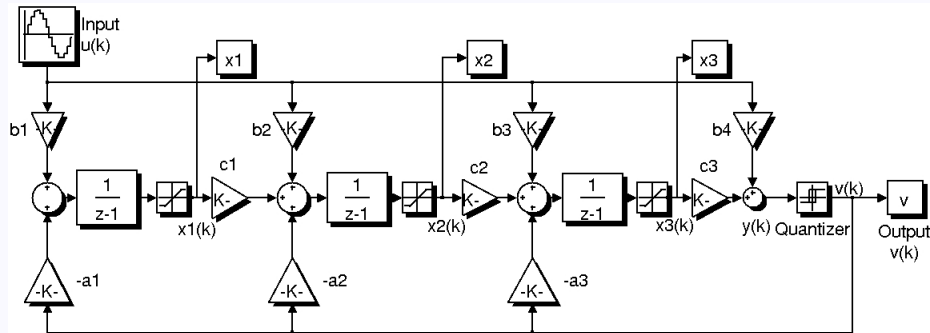
- To prove that *the system is safe over horizon N* , we prove that it is *safe at all steps $k = 1, 2, \dots, N$*

Δ - Σ modulation

- Δ - Σ modulation is a popular technique of analog-to-digital conversion
- 4 steps: anti-aliasing, oversampling, noise shaping, quantization



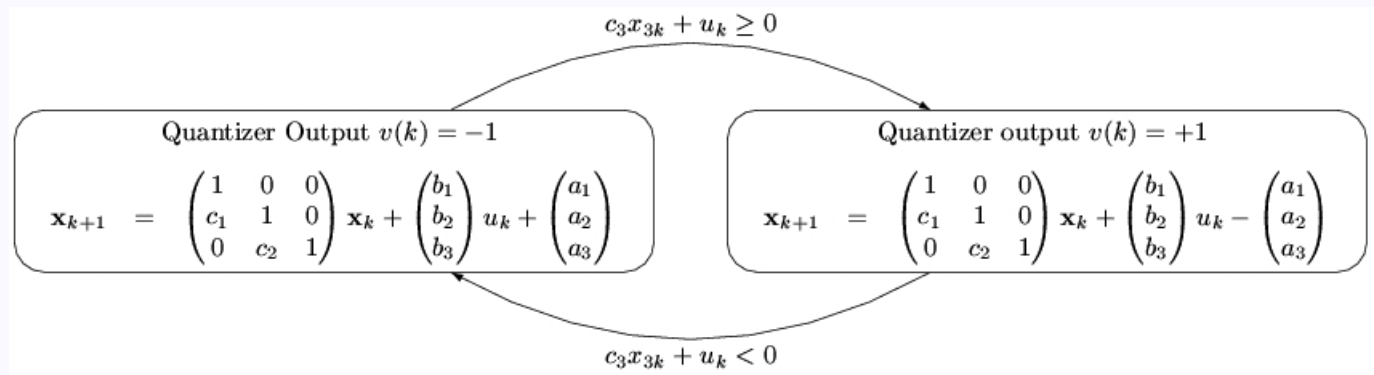
Third-order Δ - Σ modulator



- 3 integrators \Rightarrow third order Δ - Σ modulator
- High order modulators achieve better performance but induce *instability problems*
- **Stability:** a modulator is stable if the states of all its integrators always remain bounded (under bounded inputs)

Third-order Δ - Σ : hybrid model and verification

Hybrid automaton model



$\mathbf{x}_k = (x_1(k) \ x_2(k) \ x_3(k))^T$: states of the integrators

Verification

To prove that the circuit is stable, we prove that its integrators *never saturate* \Rightarrow a safety property

$$\text{Safe set } \mathcal{S} = [-x_1^{sat}, x_1^{sat}] \times [-x_2^{sat}, x_2^{sat}] \times [-x_3^{sat}, x_3^{sat}]$$

Third-order Δ - Σ verification

- Objective function can be defined as

$$J(x(k)) = \max_{i=1,2,3} (|x_i(k)| - x_i^{sat})$$

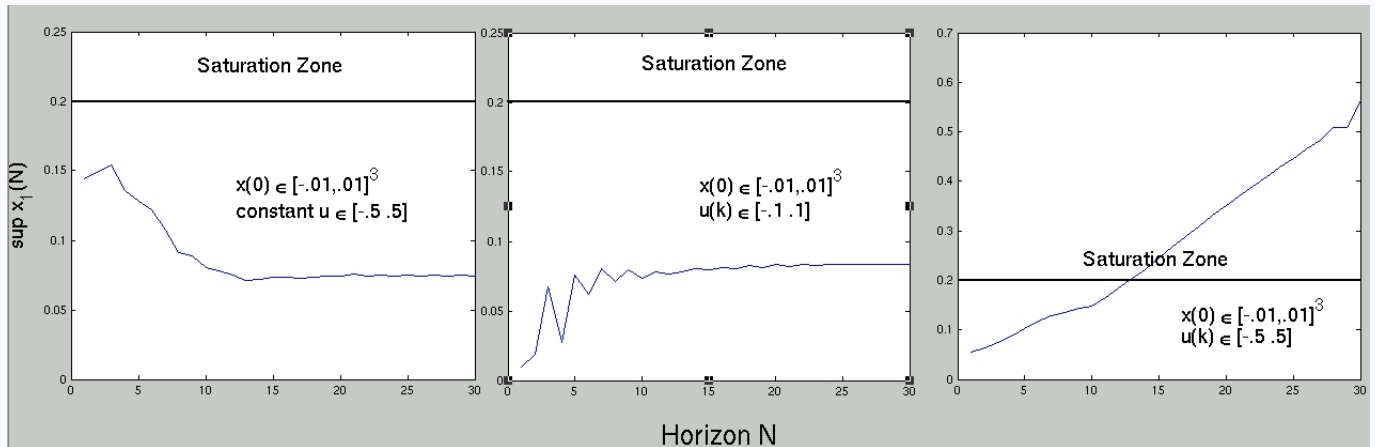
Symmetries of the dynamics and the safe set \mathcal{S} allow to split the optimization problem into three auxiliary problems:

$$\max J(x(k)) = \max_{i=1,2,3} (\max J_i(x(k)))$$

where $J_i(x(k)) = x_i(k) - x_i^{sat}$ for $i = 1, 2$, or 3 .

- Objective functions are *linear* and the dynamics of each mode is affine
 \Rightarrow *Mixed Integer Linear Programs* (for which efficient solvers exist)
- We used **MOSEK** to solve the above optimisation problem.

Δ - Σ verification: Results



Remark: frequent switching of the quantizer makes the optimisation problem very combinatorial

Conclusion

Results

- A framework for modelling and verification of analog and mixed-signal circuits using hybrid system techniques
- Encouraging experimental results

Future work

To increase scalability:

- Abstraction and model reduction techniques, exploit circuit structure
- Consider refined constraints on the input signals (e.g. bandwidth)

To integrate formal methods in the design process:

- automatic translation from circuit descriptions (such as in VHDL-AMS) to hybrid automata
- formalism for specifying properties of analog signals

End
Thank You For Your Attention