

# Verification of non-linear systems using hybridization

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# Reachability analysis

Reachable set computations are useful for

- **Verification**  
problems such as proving that the system does not reach a 'bad' state
- **Controller synthesis**  
problems such as determining the set of states from which it is possible to reach a target set while avoiding a forbidden set

Many existing methods and tools (see the next slide)

# Reachability analysis

## Direct methods

- **Track the evolution** of the reachable set under the flow of the system. Various **set representations**: *e.g.* polyhedra, ellipsoids, level sets
- **Exact** results, or **accurate approximations** with error bounds. Using **symbolic** or **numerical** computations
- **Tools**: Coho, CheckMate, d/dt, HysDel, VeriShift, Verdict, Requiem, Level-set toolbox, ..

## Indirect methods

- **Abstraction methods**: reducing to a simpler system that preserves the property (*e.g.* Tiwari & Khanna 02; Alur et al. 02; Clarke et al. 03)
  - Prove the property without computing reachable sets: *e.g.* **Barrier certificates** Prajna & Jadbabaie04, **polynomial invariants** Tiwari & Khanna04.
- ★ **Scalability** is still challenging (complexity and size of **real-life systems**)

# Plan

## Hybridization methods for non-linear systems

- Principle
- Construction of approximate systems
  - Piecewise affine approximation
  - Piecewise multi-affine approximation
- Reachability computations using approximate systems

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# Principle

System  $\Delta$  :  $\dot{x} = f(x)$ ,  $x \in \mathcal{X}$ ,  $f$  is Lipschitz

Main Idea:

- **Step 1.** Approximate the ‘complex’ system  $\Delta$  by a system which is “piecewise less complex” (complex: difficult to analyze). The approximate system is thus **hybrid**
- **Step 2.** Using the approximate system to yield approximate analysis results for  $\Delta$

Construction of the approximate system: partition the state space  $\mathcal{X}$  into disjoint regions and assign to each region an approximate vector field (which is simpler to analyze)

## Important questions :

- Error between the approximate and the original systems
- Classes of approximate vector fields to use

## Construction of the approximate system

System  $\Delta$  :  $\dot{x} = f(x)$ ,  $x \in \mathcal{X}$ ,  $f$  is Lipschitz

- Partition the state space  $\mathcal{X}$  into disjoint regions of size  $\mathbf{h}$  and assign to each region an approximate vector field
- $\mathbf{h}$ : space discretization size
- $f_{\mathbf{h}}$ : composite vector field over the whole state space  $\mathcal{X}$
- Hybridization error  $\pi(\mathbf{h}) = \sup_{x \in \mathcal{X}} \|f(x) - f_{\mathbf{h}}(x)\|$
- Conservative approximate system

System  $\Delta_{\mathbf{h}}$  :  $\dot{x} = f_{\mathbf{h}}(x) + u$

$u(\cdot)$ : disturbance taking values in  $Ball(\pi(\mathbf{h}))$

# Approximation Properties

**Conservativeness:** The set of trajectories of  $\Delta$  is included in the set of trajectories of  $\Delta_{\mathbf{h}}$

**Convergence**

- Distance between trajectories:  $x(0) = x_h(0)$ ,  $L$  is Lipschitz cst. of  $f$ .

$$\forall t \geq 0 : d(x(t), x_h(t)) \leq \frac{\pi(\mathbf{h})}{L}(e^{Lt} - 1),$$

- Distance between the reachable sets up to time  $T$

$$d(\text{Reach}_{[0,T]}(\Delta), \text{Reach}_{[0,T]}(\Delta_{\mathbf{h}})) \leq \frac{2\pi(\mathbf{h})}{L}(e^{LT} - 1),$$

that is  $\mathcal{O}(\pi(\mathbf{h}))$

- The reachable set of  $\Delta_{\mathbf{h}}$  converges to the reachable sets of  $\Delta$  with the same rate as  $f_{\mathbf{h}}$  converges to  $f$



## Preservation of attractors

If  $\Delta$  has an attractor  $A$  and  $B$  is its basin of attraction, and if

- $\forall x_0 \in B$  there exists a trajectory  $x(t)$  of  $\Delta$  such that  $x(0) = x_0$
- $\exists \delta > 0 : \mathcal{N}(A, \delta) \subseteq B$ .

Then, for all  $\varepsilon \in (0, \delta]$ , there exists  $\nu > 0$  such that if  $\pi(\mathbf{h}) \leq \nu$  then there exists a set  $A_{\mathbf{h}}$ , which is an attractor for the hybrid system  $\Delta_{\mathbf{h}}$  and such that

$$A \subseteq A_{\mathbf{h}} \subseteq \mathcal{N}(A, \varepsilon).$$

In addition,  $B$  is in the basin of attraction of  $A_{\mathbf{h}}$ .

## Approximate systems

We developed *two methods* for constructing approximate systems with good error bound  $\pi(\mathbf{h})$

- Piecewise affine systems
- Piecewise multi-affine systems

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# Piecewise affine approximation

- Using a **simplicial mesh**, each cell  $C_{\mathbf{i}}$  is a simplex of size  $\mathbf{h}$  (edge length)
- Define for each  $C_{\mathbf{i}} = \text{conv}\{v_1, \dots, v_{n+1}\}$  an **affine** function  $l_{\mathbf{i}}$  **interpolating**  $f$  at the vertices. That is, for each vertex  $v_k$  of  $C_{\mathbf{i}}$

$$l_{\mathbf{i}}(v_k) = f(v_k)$$

- Then, the approximate system is:  $f_{\mathbf{h}}(x) = l_{\mathbf{i}}(x)$  if  $x \in C_{\mathbf{i}}$ .

## Properties of the approximation:

- Given a mesh,  $f_{\mathbf{h}}$  is uniquely defined. (An affine  $n$ -dimensional function is uniquely determined by its values at  $(n+1)$  affinely independent points.)
- Piecewise affine function  $f_{\mathbf{h}}$  is **continuous over the state space**

## Hybridization error

- If  $f$  is  $L$ -Lipschitz, then

$$\pi(\mathbf{h}) \leq \mathbf{h} \frac{2n L}{n + 1},$$

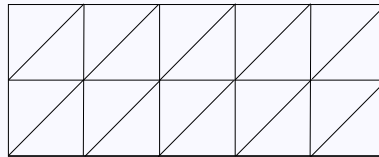
- If  $f$  is a  $C^2$  function with a second derivative bound  $K$ , then

$$\pi(\mathbf{h}) \leq \mathbf{h}^2 \frac{n^2 K}{2(n + 1)^2}.$$

$\Rightarrow$  second-order approximation  $\pi(\mathbf{h}) = \mathcal{O}(\mathbf{h}^2)$

## Mesh construction

- Define a rectangular mesh over  $\mathcal{X}$
- Decompose each hyper-rectangle into simplices: Let  $\mathcal{P}$  be the set of permutations of  $\{1, \dots, n\}$ .
  - Given  $\mathbf{p} = (\mathbf{i}_1, \dots, \mathbf{i}_n) \in \mathcal{P}$ , the set  $\mathbf{C}_{\mathbf{p}} = \{0 \leq x_{\mathbf{i}_1} \leq \dots \leq x_{\mathbf{i}_n} \leq 1\}$  is a simplex of  $\mathbb{R}^n$ .
  - $\{\mathbf{C}_{\mathbf{p}}\}_{\mathbf{p} \in \mathcal{P}}$  defines a simplicial mesh of  $[0, 1]^n$
  - Each hyper-rectangle is decomposed into  $n!$  simplices



## Analysis of approximate systems

- Reachability computations for piecewise affine system  $\Delta_h$ :  
Various existing techniques, such as **CheckMate**, **d/dt**
- Our implementation using reachability procedures of the tool **d/dt**

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# Piecewise multi-affine approximation

- Using a **rectangular mesh**, each cell  $C_i$  is a hypercube of size  $\mathbf{h}$
- Define for each cell  $C_i$  a **multi-affine** function  $f_{\mathbf{h}}$  interpolating  $f$  at its vertices
- Piecewise multi-linear function  $f_{\mathbf{h}}$  is **continuous over the state space**

**Approximation error:** If  $f$  is  $C^2$  on  $\mathcal{X}$  with bounded second order derivatives  $\Rightarrow$  **quadratic error:**  $\pi(\mathbf{h}) = \mathcal{O}(\mathbf{h}^2)$ .



## Defining multi-affine approximations

For a 2-dimensional rectangle  $[l_1, u_1] \times [l_2, u_2]$ , we define a multi-affine function that interpolates  $f$  at the vertices of the rectangle:

For  $x \in [l_1, u_1] \times [l_2, u_2]$ ,

$$\begin{aligned} f_{\mathbf{h}}(x) &= \frac{(u_2 - x_2)(u_1 - x_1)}{(u_2 - l_2)(u_1 - l_1)} f(l_1, l_2) \\ &+ \frac{(u_2 - x_2)(x_1 - l_1)}{(u_2 - l_2)(u_1 - l_1)} f(u_1, l_2) \\ &+ \frac{(x_2 - l_2)(u_1 - x_1)}{(u_2 - l_2)(u_1 - l_1)} f(l_1, u_2) \\ &+ \frac{(x_2 - l_2)(x_1 - l_1)}{(u_2 - l_2)(u_1 - l_1)} f(u_1, u_2) \end{aligned}$$

## Defining multi-affine approximations

- Iteratively applying linear interpolation on each dimension
- Let  $\mu_f(k; l_k, u_k)$  the function obtained by applying a linear interpolation of  $f$  on the  $k^{\text{th}}$  dimension at two points  $l_k$  and  $u_k$ . Let

$$\mathbf{p}(x) = \mu_f(1; l_1, u_1) \text{ and } \mathbf{q}(x) = \mu_f(2; l_2, u_2).$$

Then, for  $x \in [l_1, u_1] \times [l_2, u_2]$

$$f_{\mathbf{h}}(x) = \mu_{\mathbf{p}}(2; l_2, u_2) = \mu_{\mathbf{q}}(1; l_1, u_1).$$

# Piecewise affine vs. Piecewise multi-affine

★ Advantage comparison

<b>Simplicial meshes</b>	<b>Rectangular meshes</b>
	smaller number of cells
	less complex geometric structure
available techniques for approximate systems	???

★ Reachability computations for piecewise multi-affine systems with uncertain input:

Abstraction by projection [Asarin Dang 2004]

- Use projection to obtain a uncertain bilinear control system
- Then, use our reachability technique for bilinear control systems

Method for polynomial systems [Dang 2006]

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# Reachability computation method for polynomial systems

We consider a polynomial differential equation:

$$\dot{\mathbf{x}}(t) = g(\mathbf{x}(t))$$

Given a time step  $r$ , it is possible to find a numerical scheme (for detail see the HSCC workshop)

$$\mathbf{x}_{k+1} = Q(\mathbf{x}_k, r)$$

where  $Q(\mathbf{x}_k)$  is a polynomial in  $\mathbf{x}_k$ .

We compute this map using Bézier techniques

**The approximation error is  $\mathcal{O}(r^2)$**

# Bézier simplices - definitions

A **multi-index**  $\mathbf{i} = (\mathbf{i}[1], \dots, \mathbf{i}[n + 1])$  is a vector of non-negative integers. The norm  $\|\mathbf{i}\| = \sum_{j=1}^{n+1} \mathbf{i}[j]$

$\Delta$ :  $n$ -dimensional simplex with vertices  $\{\mathbf{v}_1, \dots, \mathbf{v}_{n+1}\}$ . Given a point  $\mathbf{x} \in \Delta$ ,  $\lambda(\mathbf{x}) = (\lambda_1(\mathbf{x}), \dots, \lambda_{n+1}(\mathbf{x}))$ : barycentric coordinates of  $\mathbf{x}$  w.r.t the vertices of  $\Delta$ , that is,  $\mathbf{x} = \sum_{k=1}^{n+1} \lambda_k(\mathbf{x})\mathbf{v}_k$  and  $\sum_{k=1}^{n+1} \lambda_k(\mathbf{x}) = 1$ .

A **Bézier simplex** of degree  $d$  is a polynomial defined as:

$$\boldsymbol{\pi}(\mathbf{x}) = \sum_{\|\mathbf{i}\|=d} \mathbf{p}_{\mathbf{i}} B_{\mathbf{i},d}(\lambda_1(\mathbf{x}), \dots, \lambda_{n+1}(\mathbf{x}))$$

where  $B_{\mathbf{i},d}(y_1, \dots, y_{n+1}) = \binom{d}{\mathbf{i}} y_1^{\mathbf{i}[1]} y_2^{\mathbf{i}[2]} \dots y_{n+1}^{\mathbf{i}[n+1]}$  (**Bernstein polynomials**), with the multimomial coefficient  $\binom{d}{\mathbf{i}} = \frac{d!}{\mathbf{i}[1]!\mathbf{i}[2]!\dots\mathbf{i}[n+1]!}$ .

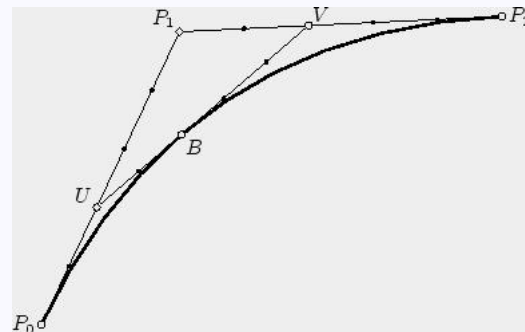
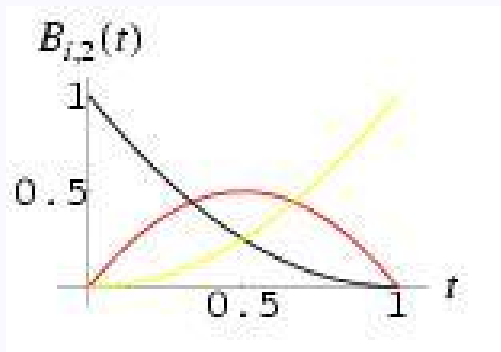
Each vector  $\mathbf{p}_{\mathbf{i}} \in \mathbb{R}^n$  is called a **Bézier control point** and all such  $\mathbf{p}_{\mathbf{i}}$  form the **Bézier control net** of  $\boldsymbol{\pi}$  with respect to  $\Delta$ .

# Example - Bézier curves

Dimension  $n = 1$ . With  $d = 2$ , multi-indices with  $||\mathbf{i}|| = 2$ :  $\{(0, 2), (1, 1), (2, 0)\}$ .

A quadratic Bézier curve is:  $\pi(\mathbf{x}) = \sum_{\mathbf{i}} \mathbf{p}_{\mathbf{i}} B_{\mathbf{i},2}(\lambda_1(\mathbf{x}), \lambda_2(\mathbf{x}))$ , for  $\mathbf{x} \in [\mathbf{v}_1, \mathbf{v}_2]$ .

There are 3 control points  $\mathbf{p}_{\mathbf{i}}$ .



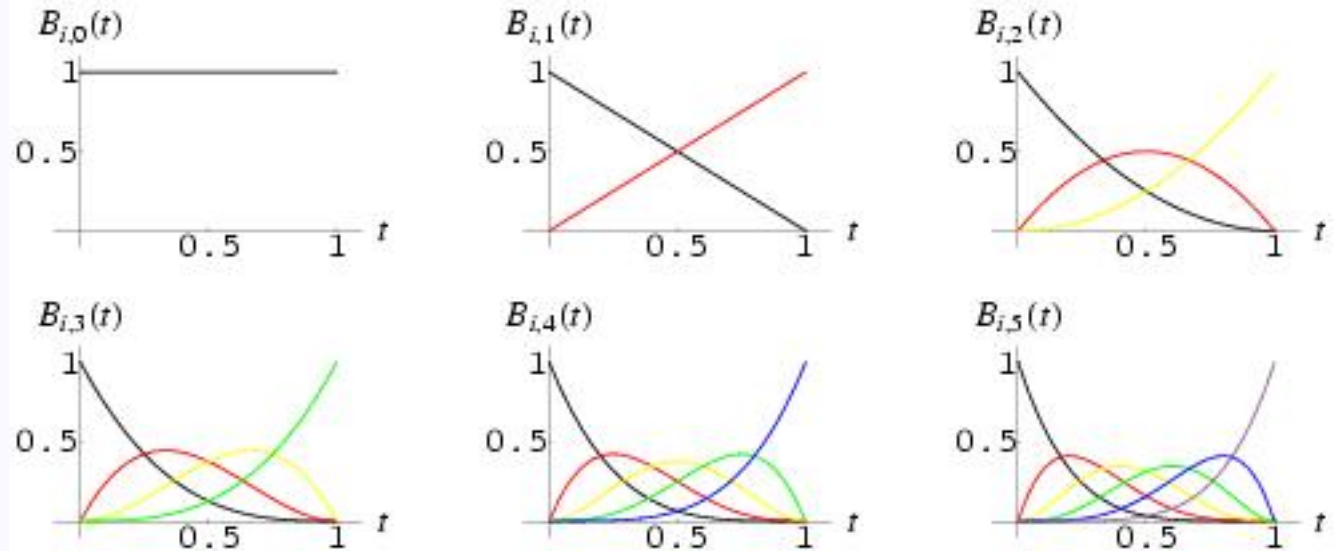
Bernstein polynomials

$$B_{\mathbf{i},2}(t) = B_{\mathbf{i},2}(y_1, y_2) = \binom{2}{\mathbf{i}} y_1^{\mathbf{i}^{[1]}} y_2^{\mathbf{i}^{[2]}}, y_1 = t \in [0, 1], y_2 = 1 - t$$



# Example - Bézier curves

Dimension  $n = 1$ .



Bernstein polynomials

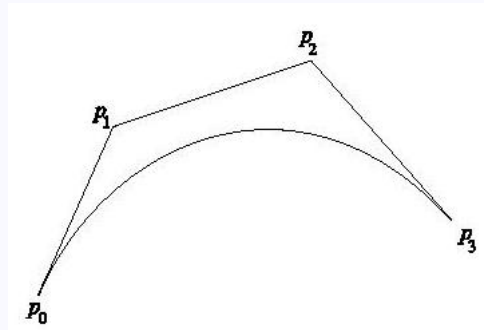
$$B_{i,d}(t) = B_{i,d}(y_1, y_2) = \binom{d}{i} y_1^{i^{[1]}} y_2^{i^{[2]}}, y_1 = t \in [0, 1], y_2 = 1 - t$$

# Bézier simplices - Shape Properties

Any polynomial can be written in form of a Bézier simplex. Given an arbitrary point  $\mathbf{x} \in \Delta$ ,

1. **Convex hull property:** the point  $\pi(\mathbf{x})$  lies inside the convex hull of the control net
2. **End-point interpolation property:** the polynomial  $\pi$  interpolates the control net at the corner control points.

Number of control points  $\mathbf{p}_i$  is  $\binom{d+n}{n} = \frac{(d+n)!}{d! n!}$ .



$\Rightarrow$  Using the convex hull of the Bézier control net as a **tight over-approximation** of  $\pi(\Delta)$ .

# Computing the Bézier control points

**Problem:** For a polynomial  $\pi$  given in monomial form and a simplex  $\Delta = \text{conv}\{\mathbf{v}_1, \dots, \mathbf{v}_{n+1}\}$ , we want to compute the control net of  $\pi$  with respect to  $\Delta$ .

**Blossoming principle** For any polynomial  $\pi : \mathbb{R}^n \rightarrow \mathbb{R}^n$  of degree  $d$ , there is a unique symmetric  $d$ -affine map  $\beta : \mathbb{R}^{nd} \rightarrow \mathbb{R}^n$  such that for all  $\mathbf{x} \in \mathbb{R}^n$ :  $p(\mathbf{x}, \dots, \mathbf{x}) = \pi(\mathbf{x})$ .

**Recall:**  $q(\mathbf{x}_1, \dots, \mathbf{x}_d)$  is  $d$ -affine if affine when all but one of its arguments are fixed; symmetric if independent of the ordering of the arguments.

Connection with the Bézier control net:

$$\mathbf{p}_i = \beta(\underbrace{\mathbf{v}_1, \dots, \mathbf{v}_1}_{i[1]}, \underbrace{\mathbf{v}_2, \dots, \mathbf{v}_2}_{i[2]}, \dots, \underbrace{\mathbf{v}_{n+1}, \dots, \mathbf{v}_{n+1}}_{i[n+1]})$$

$\Rightarrow$  To compute  $\mathbf{p}_i$  we can evaluate the blossom with some particular arguments

# Evaluating blossom values

Illustrate the computation of blossom values of polynomial  $(\mathbf{x}[i])^h(\mathbf{x}[j])^k$ .

$$\beta_{h,k}^d(\mathbf{u}_1, \dots, \mathbf{u}_d) = \frac{1}{\binom{d}{h} \binom{d-h}{k}} \sum_{\substack{I \cup J \subseteq \{1, \dots, d\}, \\ |I| = h, |J| = k, I \cap J = \emptyset}} \prod_{r \in I} \mathbf{u}_r[i] \prod_{s \in J} \mathbf{u}_s[j]$$

Denote  $\sigma_{h,k}^d = \frac{1}{\binom{d}{h} \binom{d-h}{k}} \beta_{h,k}^d(\mathbf{u}_1, \mathbf{u}_2, \dots, \mathbf{u}_d)$ .

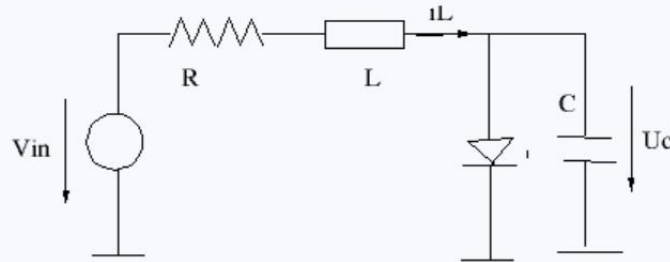
The symmetric function  $\sigma$ : choose  $h$   $i^{th}$  coordinates of  $d$  argument points and  $k$   $j^{th}$  coordinates and form their product, then sum these products over all possible choices.

$$\begin{cases} \sigma_{h,k}^d = \sigma_{h,k}^{d-1} + \mathbf{u}_d[i] \sigma_{h-1,k}^{d-1} + \mathbf{u}_d[j] \sigma_{h,k-1}^{d-1} & \text{if } h, k \geq 0 \text{ and } h+k \geq 1, \\ \sigma_{0,0}^d = 0 \end{cases}$$

Complexity  $O(d^3)$

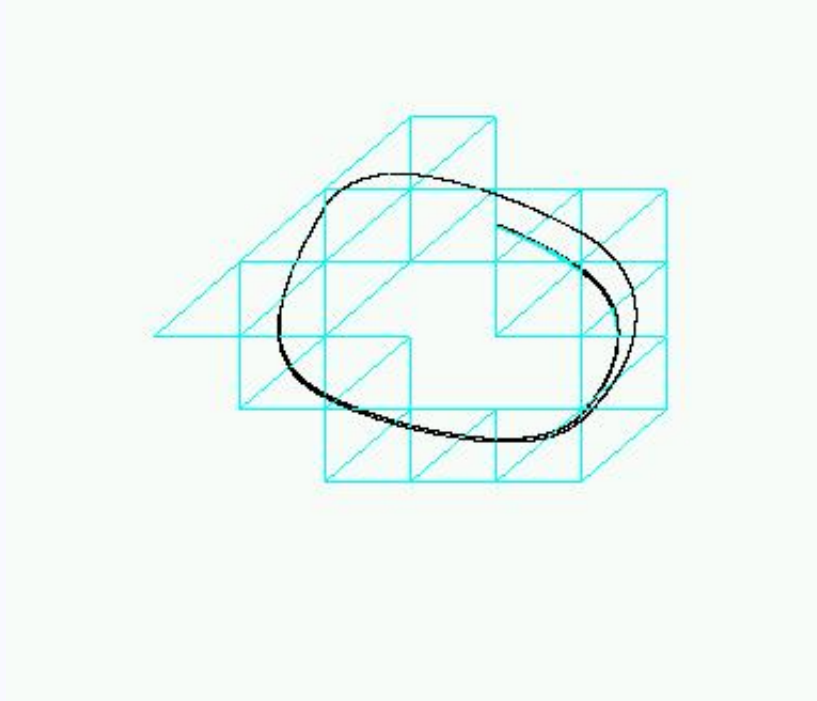


# Tunnel Diode Oscillator



$$-\frac{d}{dt}u_c = \frac{1}{C}(0.43u_c^3 - 2.69u_c^2 + 4.56u_c + 1.4(1 - \lambda)u_c) - \frac{1}{C}i_L$$
$$\frac{d}{dt}i_L = -\frac{1}{L}u_c - \frac{R}{L}i_L + \frac{1}{L}V_{in}$$

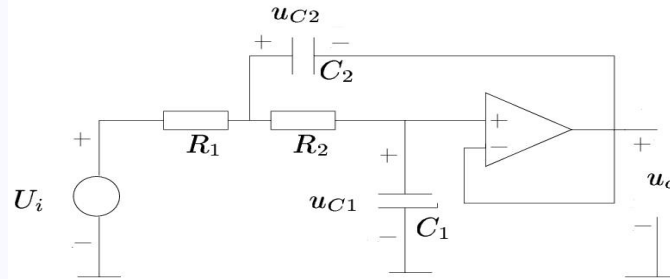
# Tunnel Diode Oscillator



The result shows a limit cycle

# Biquad low pass filter

[Hartong,Hedrich,Barke 2002]



$$\dot{u}_{C1} = \frac{u_{C2} + u_o - u_{C1}}{C_1 R_2} \quad \dot{u}_{C2} = \frac{U_i - u_{C2} - u_o}{C_2 R_1} - \frac{u_{C2} + u_o - u_{C1}}{C_2 R_2}, \quad (1)$$

$$u_o - V_{max} \tanh\left(\frac{(u_{C2} - u_o)V_e}{V_{max}}\right) + U_{om} = 0, \quad (2)$$

$$U_{om} = \mathcal{V}(i_o), \quad i_o = -C_2 \dot{u}_{C2}, \quad (3)$$

$$\mathcal{V}(i_o) = K_1 i_o + 0.5 \sqrt{K_1 i_o^2 - 2K_2 i_o I_s + K_1 I_s^2 + K_2} - 0.5 \sqrt{K_1 i_o^2 + 2K_2 i_o I_s + K_1 I_s^2 + K_2}. \quad (4)$$

Differentiating (2)  $\Rightarrow$  nonlinear ODEs on a manifold with variables  $z = (u_{C1}, u_{C2}, u_o)$



# Biquad low pass filter: Approximation by hybridization

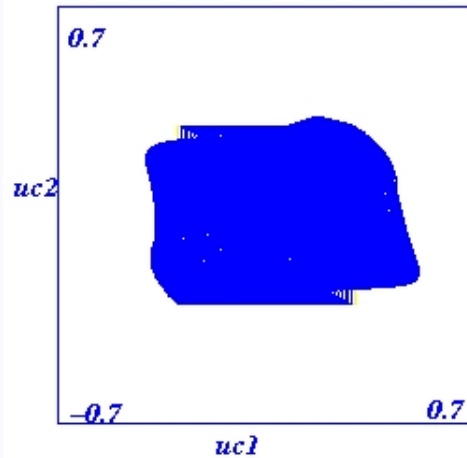
- *tanh* approximated by polynomial or piecewise affine  $\Rightarrow$  loss of important behaviors of the circuit
- the nonlinear characteristics  $U_{om} = \mathcal{V}(i_o)$  can be approximated by a piecewise-affine function:

$$U_{om} = \begin{cases} 1e10 i_o + 0.5e8 & \text{if } i_o \leq -i_s \\ 0 & \text{if } -i_s < i_o < i_s \\ 1e10 i_o - 0.5e8 & \text{if } i_o \geq i_s \end{cases} \quad (5)$$

$\Rightarrow$  Hybrid automaton with 3 continuous modes

# Biquad low pass filter: verification results

The property to verify is the *absence of overshoots*.



- $C_1 = 0.5e - 8$ ,  $C_2 = 2e - 8$ , and  $R_1 = R_2 = 1e6$  (highly damped case)
- The initial set:  $u_{C1} \in [-0.3, 0.3]$ ,  $u_{C2} \in [-0.3, 0.3]$  and  $u_o \in [-0.2, 0.2]$

# Ongoing work

- Hybridization: hierarchical mesh construction
- Using higher degree approximants, such as piecewise quadratic
- Other new classes of properties preserved by hybridization