



Measuring with Timed Patterns

CAV'15

Thomas Ferrère¹ **Oded Maler¹** **Dejan Nickovic²** **Dogan Ulus¹**

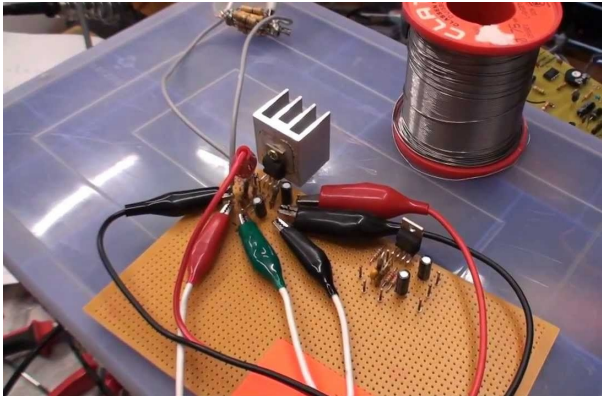
¹ VERIMAG University of Grenoble / CNRS

² AIT Austrian Institute of Technology

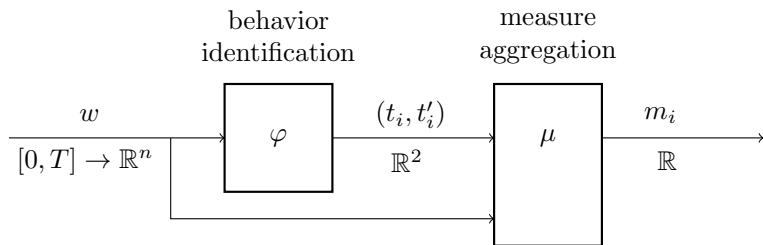
July 24, 2015

Measurements current practice...

- ▶ scripts, signal processing blocks, etc.
- ▶ ad-hoc approach

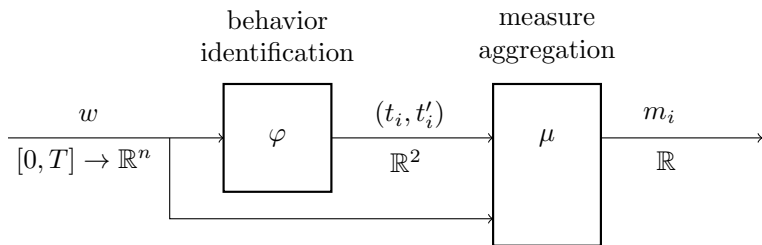


Declarative language for measurements



- ▶ **timed** regular expression φ describes intervals where measure can be taken
- ▶ continuous aggregating operators μ : duration, integral, maximum, etc.

Declarative language for measurements



- ▶ **timed** regular expression φ describes intervals where measure can be taken
- ▶ continuous aggregating operators μ : duration, integral, maximum, etc.

Timed regular expressions – interval semantics

Definition (Syntax of TRE)

$$\varphi := \epsilon \mid p \mid \varphi \cdot \varphi \mid \varphi \vee \varphi \mid \varphi \wedge \varphi \mid \varphi^* \mid \langle \varphi \rangle_{[l,u]}$$

p proposition, and l, u integer constants.

Definition (Semantics of TRE)

$$\begin{aligned}(t, t') \in [\epsilon]_w & \quad \text{iff } t = t' \\(t, t') \in [p]_w & \quad \text{iff } \forall t < t'' < t', p \in w[t''] \\(t, t') \in [\varphi \cdot \psi]_w & \quad \text{iff } \exists t \leq t'' \leq t', (t, t'') \in [\varphi]_w \text{ and } (t'', t') \in [\psi]_w \\(t, t') \in [\varphi \vee \psi]_w & \quad \text{iff } \dots \\(t, t') \in [\varphi \wedge \psi]_w & \quad \text{iff } \dots \\(t, t') \in [\varphi^*]_w & \quad \text{iff } \dots \\(t, t') \in [\langle \varphi \rangle_{[l,u]}]_w & \quad \text{iff } l \leq t' - t \leq u \text{ and } (t, t') \in [\varphi]_w\end{aligned}$$

Timed regular expressions – interval semantics

Definition (Syntax of TRE)

$$\varphi := \epsilon \mid p \mid \varphi \cdot \varphi \mid \varphi \vee \varphi \mid \varphi \wedge \varphi \mid \varphi^* \mid \langle \varphi \rangle_{[l,u]}$$

p proposition, and l, u integer constants.

Definition (Semantics of TRE)

$(t, t') \in \llbracket \epsilon \rrbracket_w$	iff	$t = t'$
$(t, t') \in \llbracket p \rrbracket_w$	iff	$\forall t < t'' < t', p \in w[t'']$
$(t, t') \in \llbracket \varphi \cdot \psi \rrbracket_w$	iff	$\exists t \leq t'' \leq t', (t, t'') \in \llbracket \varphi \rrbracket_w$ and $(t'', t') \in \llbracket \psi \rrbracket_w$
$(t, t') \in \llbracket \varphi \vee \psi \rrbracket_w$	iff	...
$(t, t') \in \llbracket \varphi \wedge \psi \rrbracket_w$	iff	...
$(t, t') \in \llbracket \varphi^* \rrbracket_w$	iff	...
$(t, t') \in \llbracket \langle \varphi \rangle_{[l,u]} \rrbracket_w$	iff	$l \leq t' - t \leq u$ and $(t, t') \in \llbracket \varphi \rrbracket_w$

Timed pattern matching

Theorem (FORMATS'14)

The set of matches $\llbracket \varphi \rrbracket_w$ is computable as a finite union of 2d zones

Proof principle

Structural induction over φ

$$\begin{aligned} z_p &\rightsquigarrow t_i < t < t' < t_{i+1} \\ z_{\varphi \cdot \psi} &\rightsquigarrow z_\varphi \circ z_\psi \\ &\dots \\ z_{(\varphi)[l, u]} &\rightsquigarrow z_\varphi \wedge l < t' - t < u \end{aligned}$$

Timed pattern matching

Theorem (FORMATS'14)

The set of matches $\llbracket \varphi \rrbracket_w$ is computable as a finite union of 2d zones

Proof principle

Structural induction over φ

$$z_p \rightsquigarrow t_i < t < t' < t_{i+1}$$

$$z_{\varphi \cdot \psi} \rightsquigarrow z_{\varphi} \circ z_{\psi}$$

...

$$z_{\langle \varphi \rangle_{[l, u]}} \rightsquigarrow z_{\varphi} \wedge l < t' - t < u$$

Example

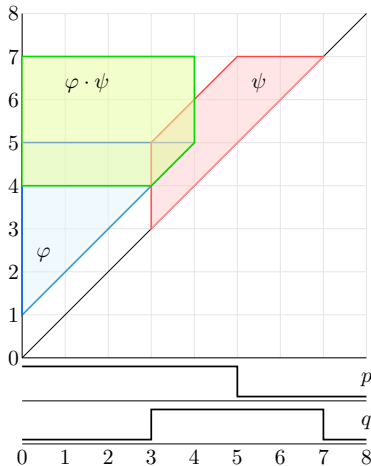
Expressions:

$$\varphi = \langle p \rangle_{[1,5]}$$

$$\psi = \langle q \rangle_{[0,2]}$$

$$\varphi \cdot \psi$$

Set of matches:



Example

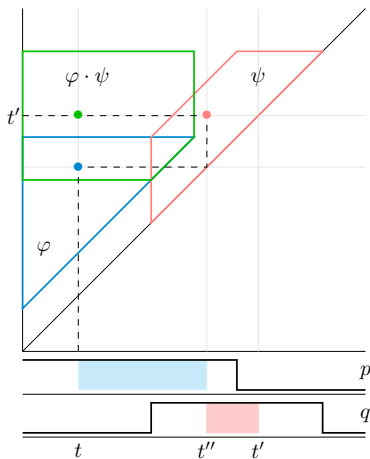
Expressions:

$$\varphi = \langle p \rangle_{[1,5]}$$

$$\psi = \langle q \rangle_{[0,2]}$$

$$\varphi \cdot \psi$$

Set of matches:



Conditional expressions

Introduce **preconditions** and **postconditions**.

Definition (Syntax of Conditional TRE)

$$\varphi := \dots \mid \varphi \cdot \varphi \mid \dots \mid \varphi ? \varphi \mid \varphi ! \varphi$$

Definition (Semantics of Conditional TRE)

...

$$(t, t') \in \llbracket \varphi \cdot \psi \rrbracket_w \quad \text{iff} \quad \exists t \leq t'' \leq t' \quad (t, t'') \in \llbracket \varphi \rrbracket_w \quad \text{and} \quad (t'', t') \in \llbracket \psi \rrbracket_w$$

...

$$(t, t') \in \llbracket \varphi ? \psi \rrbracket_w \quad \text{iff} \quad \exists t'' \leq t \quad (t, t'') \in \llbracket \varphi \rrbracket_w \quad \text{and} \quad (t'', t) \in \llbracket \psi \rrbracket_w$$

$$(t, t') \in \llbracket \varphi ! \psi \rrbracket_w \quad \text{iff} \quad \exists t' \leq t'' \quad (t, t') \in \llbracket \varphi \rrbracket_w \quad \text{and} \quad (t', t'') \in \llbracket \psi \rrbracket_w$$

Conditional expressions

Introduce **preconditions** and **postconditions**.

Definition (Syntax of Conditional TRE)

$$\varphi := \dots \mid \varphi \cdot \varphi \mid \dots \mid \varphi ? \varphi \mid \varphi ! \varphi$$

Definition (Semantics of Conditional TRE)

$$\begin{aligned} & \dots \\ (t, t') \in \llbracket \varphi \cdot \psi \rrbracket_w & \text{ iff } \exists t \leq t'' \leq t' \quad (t, t'') \in \llbracket \varphi \rrbracket_w \quad \text{and} \quad (t'', t') \in \llbracket \psi \rrbracket_w \\ & \dots \\ (t, t') \in \llbracket \psi ? \varphi \rrbracket_w & \text{ iff } \exists t'' \leq t \quad (t, t') \in \llbracket \varphi \rrbracket_w \quad \text{and} \quad (t'', t) \in \llbracket \psi \rrbracket_w \\ (t, t') \in \llbracket \varphi ! \psi \rrbracket_w & \text{ iff } \exists t' \leq t'' \quad (t, t') \in \llbracket \varphi \rrbracket_w \quad \text{and} \quad (t', t'') \in \llbracket \psi \rrbracket_w \end{aligned}$$

Conditional expressions

Introduce **preconditions** and **postconditions**.

Definition (Syntax of Conditional TRE)

$$\varphi ::= \dots \mid \varphi \cdot \varphi \mid \dots \mid \varphi ? \varphi \mid \varphi ! \varphi$$

Definition (Semantics of Conditional TRE)

$$\begin{array}{l} \dots \\ (t, t') \in \llbracket \psi \cdot \psi \rrbracket_w \quad \text{iff} \quad \exists t \leq t'' \leq t' \quad (t, t'') \in \llbracket \varphi \rrbracket_w \quad \text{and} \quad (t'', t') \in \llbracket \psi \rrbracket_w \\ \dots \\ (t, t') \in \llbracket \psi ? \varphi \rrbracket_w \quad \text{iff} \quad \exists t'' \leq t \quad (t, t'') \in \llbracket \varphi \rrbracket_w \quad \text{and} \quad (t'', t) \in \llbracket \psi \rrbracket_w \\ (t, t') \in \llbracket \varphi ! \psi \rrbracket_w \quad \text{iff} \quad \exists t' \leq t'' \quad (t, t') \in \llbracket \varphi \rrbracket_w \quad \text{and} \quad (t', t'') \in \llbracket \psi \rrbracket_w \end{array}$$

Example

Expressions:

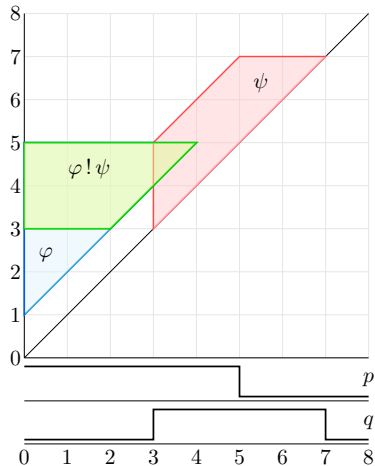
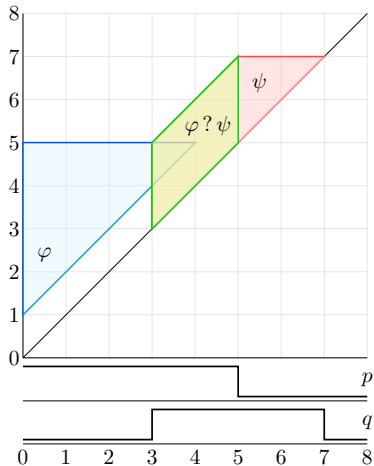
$$\varphi = \langle p \rangle_{[1,5]}$$

$$\psi = \langle q \rangle_{[0,2]}$$

$$\varphi! \psi$$

$$\varphi! \psi$$

Set of matches:



Expressions with events

Events

Zero-duration expressions:

$$\begin{array}{ll} \uparrow p = \bar{p} ? \epsilon ! p & \text{(rising edge)} \\ \downarrow p = p ? \epsilon ! \bar{p} & \text{(falling edge)} \end{array}$$

Event-bounded expressions

Syntactically enforced:

$$\psi := \uparrow p \mid \downarrow p \mid \psi \cdot \varphi \cdot \psi \mid \psi \cup \psi \mid \psi \cap \varphi$$

with φ arbitrary expression

Proposition (Finiteness)

Event-bounded expressions have a finite set of matches.

Expressions with events

Events

Zero-duration expressions:

$$\begin{array}{ll} \uparrow p = \bar{p} ? \epsilon ! p & \text{(rising edge)} \\ \downarrow p = p ? \epsilon ! \bar{p} & \text{(falling edge)} \end{array}$$

Event-bounded expressions

Syntactically enforced:

$$\psi := \uparrow p \mid \downarrow p \mid \psi \cdot \varphi \cdot \psi \mid \psi \cup \psi \mid \psi \cap \varphi$$

with φ arbitrary expression

Proposition (Finiteness)

Event-bounded expressions have a finite set of matches.

Expressions with events

Events

Zero-duration expressions:

$\uparrow p = \bar{p} ? \epsilon ! p$ (rising edge)

$\downarrow p = p ? \epsilon ! \bar{p}$ (falling edge)

Event-bounded expressions

Syntactically enforced:

$\psi := \uparrow p \mid \downarrow p \mid \psi \cdot \varphi \cdot \psi \mid \psi \cup \psi \mid \psi \cap \varphi$

with φ arbitrary expression

Proposition (Finiteness)

*Event-bounded expressions have a **finite** set of matches.*

Example

Expressions:

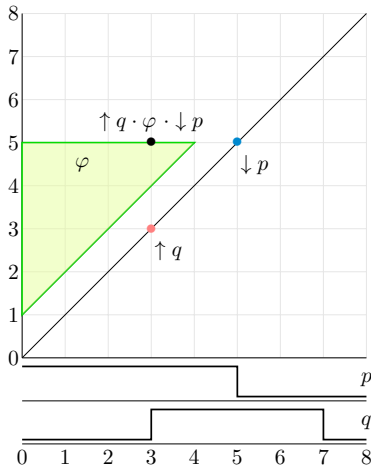
$\downarrow p$

$\uparrow q$

$\varphi = \langle p \rangle_{[1,5]}$

$\uparrow q \cdot \varphi \cdot \downarrow p$

Set of matches:



Measurements

Measure Pattern

A Conditional TRE

$$\varphi = \alpha ? \psi ! \beta$$

with arbitrary conditions α , β , and ψ event-bounded.

Measure Expression

An expression

$$\mu(\varphi)$$

with φ a measure pattern, and $\mu = \text{duration}, \text{sup}_{x}, \text{integral}_{x}, \dots$
continuous aggregation operator.

Measurements

Measure Pattern

A Conditional TRE

$$\varphi = \alpha ? \psi ! \beta$$

with arbitrary conditions α , β , and ψ event-bounded.

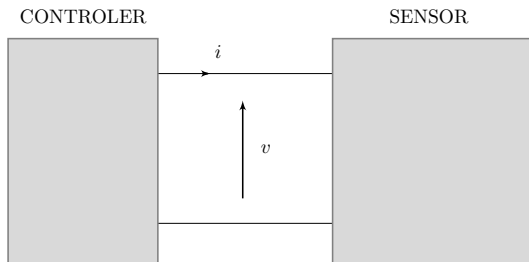
Measure Expression

An expression

$$\mu(\varphi)$$

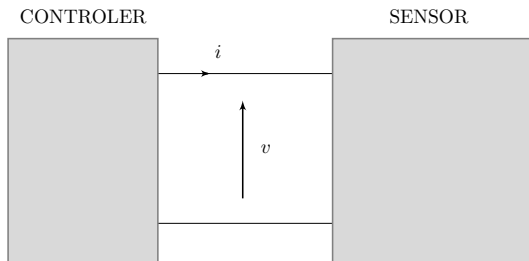
with φ a measure pattern, and $\mu =$ **duration**, **sup_x**, **integral_x**, ...
continuous aggregation operator.

DSI3 standard



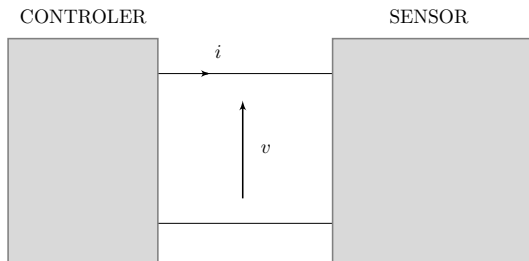
- ▶ Analog communication protocol
- ▶ Communication via **pulses** on
 - ▶ voltage line v
 - ▶ current line i
- ▶ Two phases with different **nominal levels**
 - ▶ discovery mode: v in range V_0 to V_1
 - ▶ command and response mode: v in range V_2 to V_3

DSI3 standard



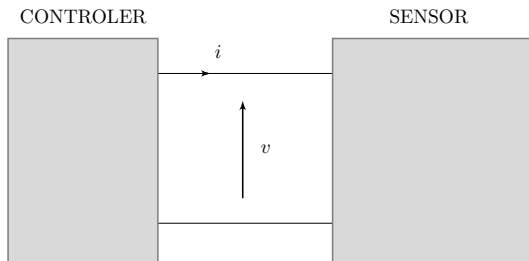
- ▶ Analog communication protocol
- ▶ Communication via **pulses** on
 - ▶ voltage line v
 - ▶ current line i
- ▶ Two phases with different **nominal levels**
 - ▶ discovery mode: v in range V_0 to V_1
 - ▶ command and response mode: v in range V_2 to V_3

DSI3 standard



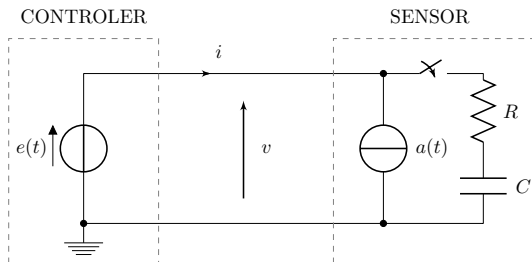
- ▶ Analog communication protocol
- ▶ Communication via **pulses** on
 - ▶ voltage line v
 - ▶ current line i
- ▶ Two phases with different **nominal levels**
 - ▶ discovery mode: v in range V_0 to V_1
 - ▶ command and response mode: v in range V_2 to V_3

Model and requirements



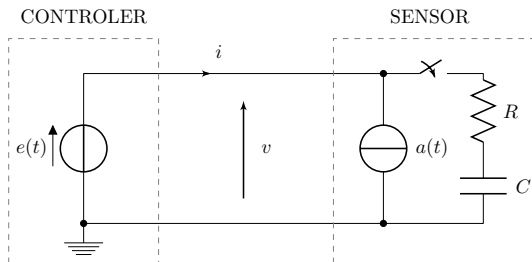
- ▶ Behavioral model:
 - ▶ gaussian distribution of pulse timing
 - ▶ uniform distribution of sensor load resistance
- ▶ Simulation: 100 sequences of discovery + command and response
- ▶ Measurements:
 1. time between consecutive discovery pulses
 2. energy transmitted through power pulses

Model and requirements



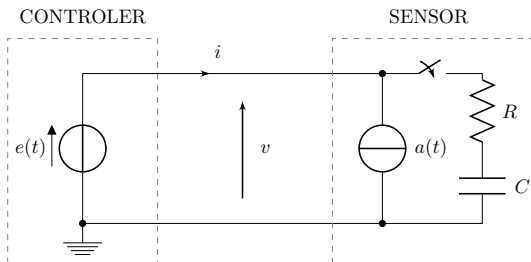
- ▶ Behavioral model:
 - ▶ gaussian distribution of pulse timing
 - ▶ uniform distribution of sensor load resistance
- ▶ Simulation: 100 sequences of discovery + command and response
- ▶ Measurements:
 1. time between consecutive discovery pulses
 2. energy transmitted through power pulses

Model and requirements



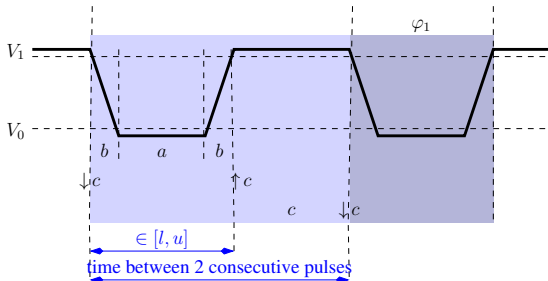
- ▶ Behavioral model:
 - ▶ gaussian distribution of pulse timing
 - ▶ uniform distribution of sensor load resistance
- ▶ Simulation: 100 sequences of discovery + command and response
- ▶ Measurements:
 1. time between consecutive discovery pulses
 2. energy transmitted through power pulses

Model and requirements



- ▶ Behavioral model:
 - ▶ gaussian distribution of pulse timing
 - ▶ uniform distribution of sensor load resistance
- ▶ Simulation: 100 sequences of discovery + command and response
- ▶ Measurements:
 1. time between consecutive discovery pulses
 2. energy transmitted through power pulses

Measurement 1: time between consecutive discovery pulses



- ▶ Voltage levels:

$$a \equiv v \leq V_0$$

$$b \equiv V_0 \leq v \leq V_1$$

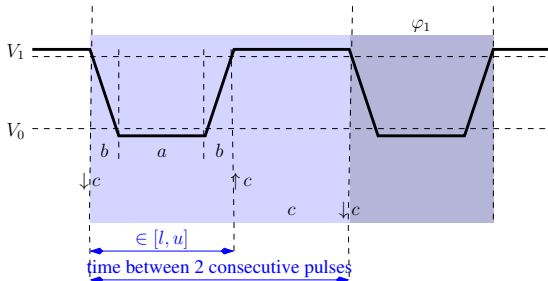
$$c \equiv v \geq V_1$$

- ▶ Pulse pattern:

$$\varphi_1 \equiv \downarrow c \cdot \langle b \cdot a \cdot b \rangle_{[l, u]} \cdot \uparrow c$$

- ▶ Measure expression: $M_1 = \mathbf{duration}(\varphi_1 \cdot c! \varphi_1)$

Measurement 1: time between consecutive discovery pulses



- ▶ Voltage levels:

$$a \equiv v \leq V_0$$

$$b \equiv V_0 \leq v \leq V_1$$

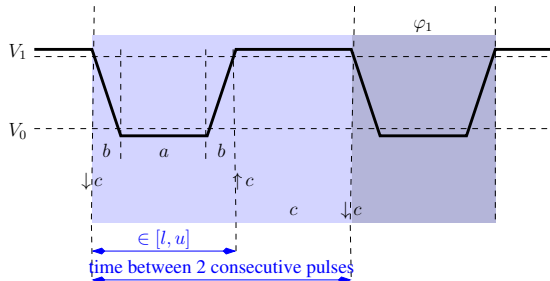
$$c \equiv v \geq V_1$$

- ▶ Pulse pattern:

$$\varphi_1 \equiv \downarrow c \cdot \langle b \cdot a \cdot b \rangle_{[l, u]} \cdot \uparrow c$$

- ▶ Measure expression: $M_1 = \mathbf{duration}(\varphi_1 \cdot c! \varphi_1)$

Measurement 1: time between consecutive discovery pulses



- ▶ Voltage levels:

$$a \equiv v \leq V_0$$

$$b \equiv V_0 \leq v \leq V_1$$

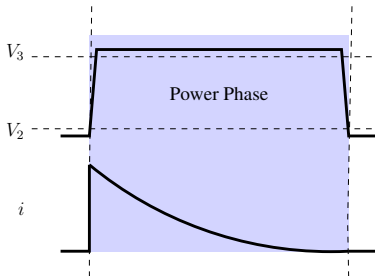
$$c \equiv v \geq V_1$$

- ▶ Pulse pattern:

$$\varphi_1 \equiv \downarrow c \cdot \langle b \cdot a \cdot b \rangle_{[l, u]} \cdot \uparrow c$$

- ▶ Measure expression: $M_1 = \mathbf{duration}(\varphi_1 \cdot c! \varphi_1)$

Measurement 2: energy transmitted during power pulses



- ▶ Voltage levels:

$$e \equiv v \geq V_2$$

$$f \equiv V_2 \leq v \leq V_3$$

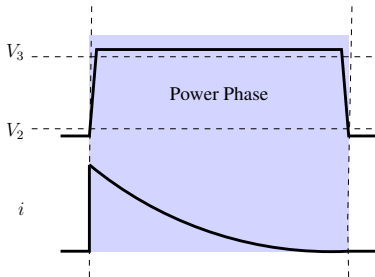
$$g \equiv v \geq V_3$$

- ▶ Pulse pattern:

$$\varphi_2 \equiv \uparrow e \cdot f \cdot g \cdot f \cdot \downarrow e$$

- ▶ Measure expression: $M_2 := \text{integral}_{v \times i}(\varphi_2)$

Measurement 2: energy transmitted during power pulses



- ▶ Voltage levels:

$$e \equiv v \geq V_2$$

$$f \equiv V_2 \leq v \leq V_3$$

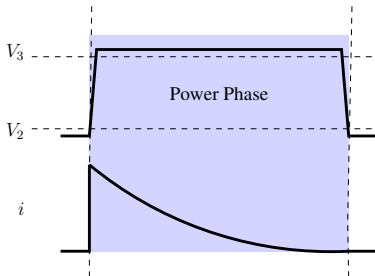
$$g \equiv v \geq V_3$$

- ▶ Pulse pattern:

$$\varphi_2 \equiv \uparrow e \cdot f \cdot g \cdot f \cdot \downarrow e$$

- ▶ Measure expression: $M_2 := \text{integral}_{v \times i}(\varphi_2)$

Measurement 2: energy transmitted during power pulses



- ▶ Voltage levels:

$$e \equiv v \geq V_2$$

$$f \equiv V_2 \leq v \leq V_3$$

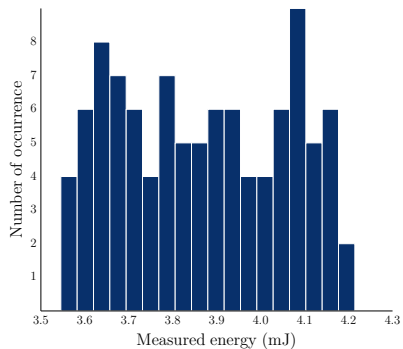
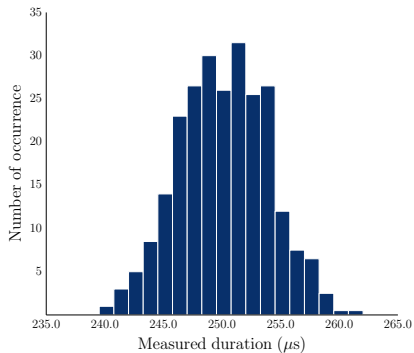
$$g \equiv v \geq V_3$$

- ▶ Pulse pattern:

$$\varphi_2 \equiv \uparrow e \cdot f \cdot g \cdot f \cdot \downarrow e$$

- ▶ Measure expression: $M_2 := \mathbf{integral}_{v \times i}(\varphi_2)$

Results



Performance

Computation time (seconds) relative to sampling rate:

samples	Measure 1				Measure 2			
	T_p	T_φ	T_μ	T	T_p	T_φ	T_μ	T
1M	0.047	0.617	0.000	0.664	0.009	0.004	0.011	0.024
5M	0.197	0.612	0.000	0.809	0.050	0.005	0.047	0.103
10M	0.386	0.606	0.000	0.992	0.101	0.005	0.100	0.216
20M	0.759	0.609	0.000	1.368	0.203	0.005	0.260	0.468

Program:

- ▶ TRE matching algorithms based on IF library
- ▶ Python signal processing library

Conclusion

Present

- ▶ declarative language for mixed-signal measurements
- ▶ general and efficient to monitor

Future

- ▶ language extension
- ▶ online measurements