

The Coq proof assistant : principles and practice

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Lecture 7

Polymorphism

Polymorphism

Lists

Polymorphism

Lists

A type can be a parameter of a function

Example: the identity function

Definition `ide := fun (X: Type) => fun (x: X) => x.`

Definition `ide (X: Type) (x: X) := x.`

Implicit arguments

When using the identity function, the first argument can be automatically inferred from the second

Example

```
id nat 3
```

```
id _ 3
```

Local declaration

```
Definition id {X: Type} (x: X) := x.
```

Simplified application

```
id 3
```

Implicit arguments

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id nat 3
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Local declaration

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Definition id {X: Type} (x: X) := x.
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Simplified application

```
id 3
```

Recovering explicit application

```
@id nat
```

```
id (X:=nat)
```

Implicit arguments

When using the identity function, the first argument can be automatically inferred from the second

Example

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id nat 3
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Local declaration

```
Definition id {X: Type} (x: X) := x.
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Simplified application

```
id 3
```

Recovering explicit application

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@id nat
```

```
id (X:=nat)
```

Global declaration

```
Set Implicit Arguments.
```


Polymorphism

Lists

Polymorphic inductive definition

```
Inductive list (X: Set) : Set :=  
  | nil : list X  
  | cons : X -> list X -> list X.
```

On Type

Can be used in more situations (e.g., lists of predicates)

```
Inductive list (X: Type) : Type :=  
  | nil : list X  
  | cons : X -> list X -> list X.
```

Basic important properties

Coq

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Polymorphism

Lists

`app` : for appending two lists

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`nil` is **neutral** on the left and on the right for `app`

- ▶ **left** : by reflexivity
- ▶ **right** : by induction

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- ▶ **right** : by induction

`app` is **associative**

- ▶ `app (app u v) w = app u (app v w)`
just by induction on **u**

See coq files Lecture07_lists