

A LUSTRE V6 TUTORIAL

Verimag

December 5, 2008 -

Outline

- **Lustre**
- **Lustre V6**
- **The Lustre V6 compiler**

Outline

- Lustre

- Lustre V6

P. Raymond & the Synchronous group et al.

- The Lustre V6 compiler

P. Raymond, J. Ballet, E. Jahier

Lustre

a Data-flow Synchronous Language

- Generalised synchronous circuits: wires hold numerics
- Operators + wires structured into nodes
- Pre-defined operators
 - Boolean: and, not, ...
 - Arithmetic: +, -, ...
 - Temporal: pre, when, current

Lustre

Targetting reactive critical systems

- **Time constraints**

 - we want a predictable bound on execution time

- **Memory constraints**

 - we want a predictable bound on memory usage

 - (we want that bound to be as small as possible)

 - ⇒ **No loops, first-order**

Lustre

a loop-free first-order language

But Can those limitations be overlooked ? ■

→ Yes: loops and genericity were introduced in V4

Lustre

Example of loops and genericity in V4

```
node add(const n:int; t1,t2 : int ^ n)
returns (res:int ^ n);
let
  res = t1 + t2; -- for i=0..n-1, res[i] = t1[i] + t2[i];
el
```

- this is legal as long as n is a ground constant which value is known at **compile time** \rightarrow static genericity
- Pushing that idea further \Rightarrow Lustre V6

Outline

- Lustre

- Lustre V6

a statically generic (1.5-order) Lustre

- The Lustre V6 compiler

Lustre V6

What's new (compared to V4)

- **Structure and enumerated types**
- **Package mechanism (Ada-like)**
 - Name space
 - Encapsulation
- **(Static) Genericity**
 - Parametric packages
 - Parametric nodes (well-typed macros)
 - Static recursion
 - Array iterators (versus homomorphic extension – not new; different)

Lustre V6

Structures

```
type complex = struct {  
    re : real = 0.;  
    im : real = 0.  
};
```

```
node plus (a, b : complex) returns (c : complex);  
let  
    c = complex { re = a.re+b.re ; im = a.im+b.im };  
tel
```

Lustre V6

Enumerated type

```
type trival = enum { Pile, Face, Tranche };
```

Lustre V6

Enumerated clocks + merge (©Pouzet)

```
type trival = enum { Pile, Face, Tranche };
node merge_node(clk: trival;
    i1 when Pile(clk); i2 when Face(clk);
    i3 when Tranche(clk))
returns (y: int);
let
    y = merge clk
        (Pile:    i1)
        (Face:    i2)
        (Tranche: i3);
tel
```

Lustre V6

Packages

```
package complex
provides
```

```
  type t; -- Encapsulation
```

```
  const i:t;
```

```
  node re(c: t) returns (r:real);
```

```
body
```

```
  type t = struct { re : real ; im : real };
```

```
  const i:t = t { re = 0. ; im = 1. };
```

```
  node re(c: t) returns (re:real);
```

```
  let re = c.re; tel;
```

```
end
```

Lustre V6

Generic packages

```
model modSimple
  needs type t;
  provides
    node fby1(init, fb: t) returns (next: t);
body
  node fby1(init, fb: t) returns (next: t);
  let next = init -> pre fb; tel
end
package pint is modSimple(t=int);
```

Lustre V6

Generic nodes

```
node mk_tab<<type t; const init: t; const size: int>>  
    (a:t) returns (res: t^size);  
let  
    res =  init ^ size;  
tel █  
node tab_int3 = mk_tab<<int, 0, 3>>; █  
node tab_bool4 = mk_tab<<bool, true, 4>>;
```


Lustre V6

Generic nodes

```
node toto_n<<
  node f(a, b: int) returns (x: int);
  const n : int
  >>(a: int) returns (x: int^n);
var v : int;
let
  v = f(a, 1);
  x = v ^ n;
tel
node toto_3 = toto_n<<Lustre::iplus, 3>>;
```

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Static recursion

```
node consensus<<const n : int>>(T: bool^n)
returns (a: bool);
let
    a = with (n = 1) then T[0]
        else T[0] and consensus << n-1 >> (T[1 .. n-1]);
tel

node main = consensus<<8>>;
```

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Are parametric nodes necessary?

- Indeed, parametric nodes could be emulated with the package mechanism

→ but we keep them to keep the syntax light

→ we didn't really want to have recursive packages

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Arrays

- **As in Lustre V4**

- **The array size is static** (`var mat23: int ^ 2 ^ 3;`)

- **Array slices** (`T1[3..5] = T2[0..2];`)

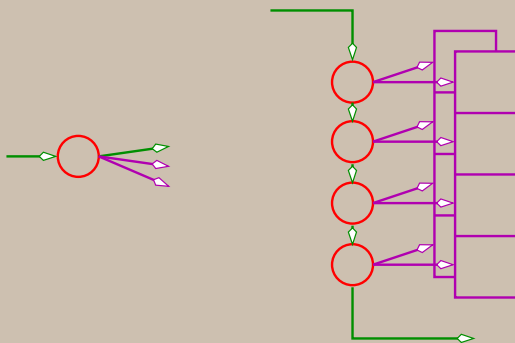
- **But no more homomorphic extension**

- where $t1 + t2$ means $\forall i \in \{0, \dots, size - 1\}, t1[i] + t2[i]$

- ⇒ **operate on arrays via iterators**

Lustre V6

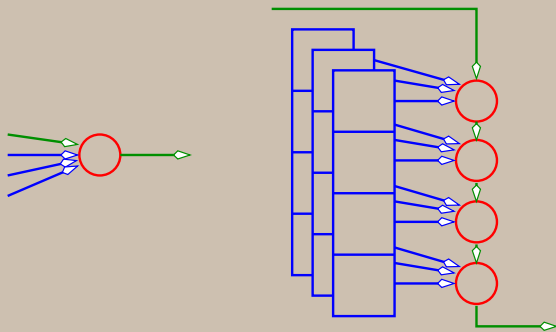
The fill iterator



```
node incr (acc : int) returns (acc', res : int);  
fill<<incr; 4>>(0)  $\rightsquigarrow$  (4, [0,1,2,3])
```

Lustre V6

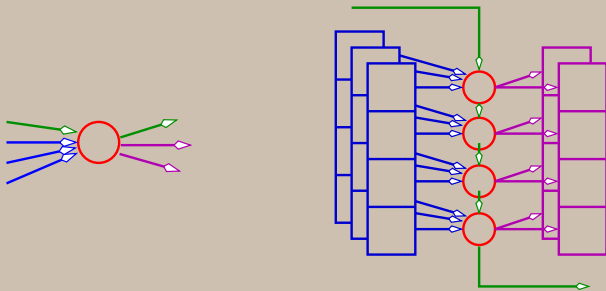
The red iterator



| `red<<+; 3>>(0, [1,2,3]) \rightsquigarrow 6`

Lustre V6

fill+red=mapred, fillred, fold

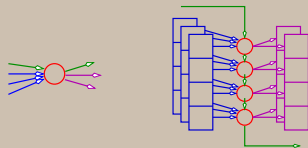


fill<<incr; 4>>(0) \equiv **fold**<<incr; 4>>(0)
red<<+; 3>>(0, [1,2,3]) \equiv **fold**<<+; 3>>(0, [1,2,3])

Lustre V6

The fold iterator

```
node cumul(acc_in,x:int) returns (acc_out,y:int)
let
  y = acc_in+x;
  acc_out = y;
tel
```

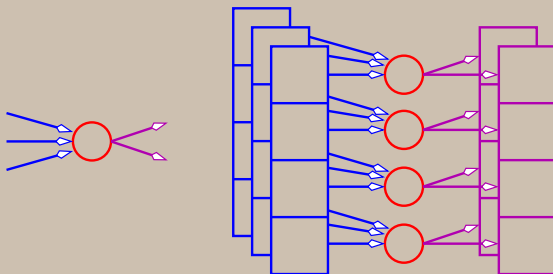


```
fold<<cumul>>(0, [1,2,3])  $\rightsquigarrow$  (6, [1,3,6])
```

```
fold<<fold<<fold<<full_adder; n>>; m>>; p>>
  (false, x, y)  $\rightsquigarrow$  (r, ''x+y'')
```


Lustre V6

The map iterator



`map <<+; 3>>([1,0,2], [3,6,-1]) \rightsquigarrow [4,6,1]`

Lustre V6

About Lustre V6 array iterators

More general than usual iterators:

they are of variable arity

Outline

- **Lustre**
- **Lustre V6**
- **The Lustre V6 compiler**
 - **The front-end**
 - **The back-end (J. Ballet)**
 - **The back-back-end (J. Ballet)**

The Lustre V6 compiler

The Front-end: LUS2LIC

- **Perform usual checks**
 - Syntax, Types, Clocks
 - Unique definition of outputs
 - Combinational cycles detection
- **Perform some static evaluation**
 - arrays size
 - parametric packages and nodes
 - recursive nodes
- **Generate intermediate code: LIC (Lustre internal code)**

Lustre Internal Code (LIC)

was: expanded code (ec)

- LIC \equiv core Lustre
 - No more packages
 - Parametric constructs are instantiated
 - constants
 - types
 - nodes

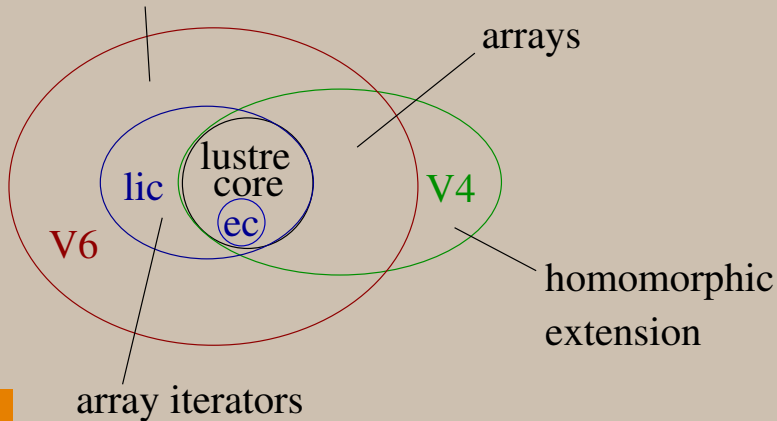
Lustre Internal Code (LIC)

was: expanded code (ec) cont.

- LIC versus ec
 - Nodes are not (necessarily) expanded
 - Arrays are not (necessarily) expanded
- LIC versus Lustre v4
 - Structures and enums
 - array iterators

Lustre potatoes

struct, enums, packages, genericity, ...



The Lustre V6 compiler

The back-end

The role of the backend is to generate sequential code

We defined (yet) another intermediary format to represent sequential code: SOC (Synchronous Object Code)

The idea is that translating this format into any sequential language is easy, and done at the very end

The back-end

maps each node to a Synchronous Object Component (SOC)

- A SOC is made of:
 - a set of memories
 - a set of methods: typically, an init and a step method
- each method is made of a sequence of guarded atomic operations
- atomic operation (named actions) can be
 - another SOC method call
 - an assignment (a wire)

The back-end

From node to SOC

For each node, we:

- Identify memories
- Explicitely separate the control (clocks) from the computations
 - set of guarded equations
- Split equations into more finer-grained steps: actions
 - a set of guarded actions (a wire or a call)
- Find a correct ordering for actions (sheduling)
 - a sequence of guarded actions

The back-back-end

From SOC to C

- pretty-print the SOC into, let's say, C
- provide a C implementation of every predefined (non-temporal) operators

Lustre V6 compiler

An alpha release is available

<http://www-verimag.imag.fr/~synchron/lustre-v6/>

- **The front-end `lus2lic` seems ok**
- **`lus2lic --lustre-v4`: added last friday; seems to work**
- **The back-back: generates C code... But its not finished.**

Thanks for your attention